

# Navigating Internet Research with P4: Solutions for Performance and Security

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[netsyn.princeton.edu](https://netsyn.princeton.edu)





What do we want from the Internet?

# What do we want from the Internet?

Cyber-physical systems

Live streaming

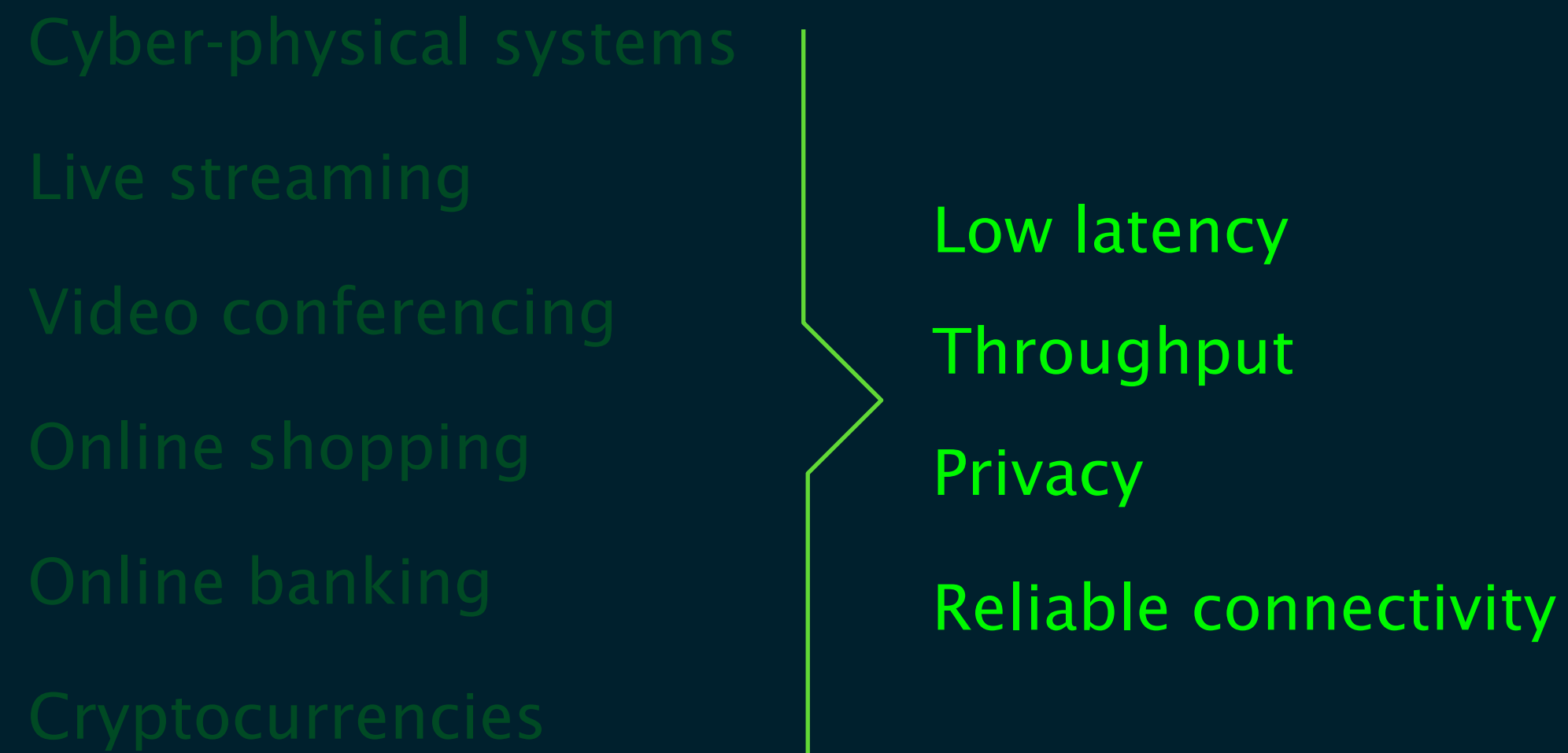
Video conferencing

Online shopping

Online banking

Cryptocurrencies

# What do we want from the Internet?



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Cyber-physical systems

Today's Internet provides best-effort service

Cryptocurrencies

...leading to performance, privacy, and security problems







Internet research is hindered by both protocols and hardware





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## BGP....

- lack of route control
- suboptimal routing
- insecure routing
- lack of path diversity

...





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## BGP....

- lack of route control
- suboptimal routing
- insecure routing
- lack of path diversity

...



## Internet Routers....

- fixed-headers support
- no cryptographic operation
- lack of performance visibility
- no DDoS support

....



What can you do with a couple of programmable points in the Internet?



What can you do with a couple of programmable points in the Internet?

Tango: performance-driven  
routing system

NSDI'24

SABRE: secure overlay  
for BTC block propagation

NDSS'19



What can you do with a couple of programmable points in the Internet?

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routing system

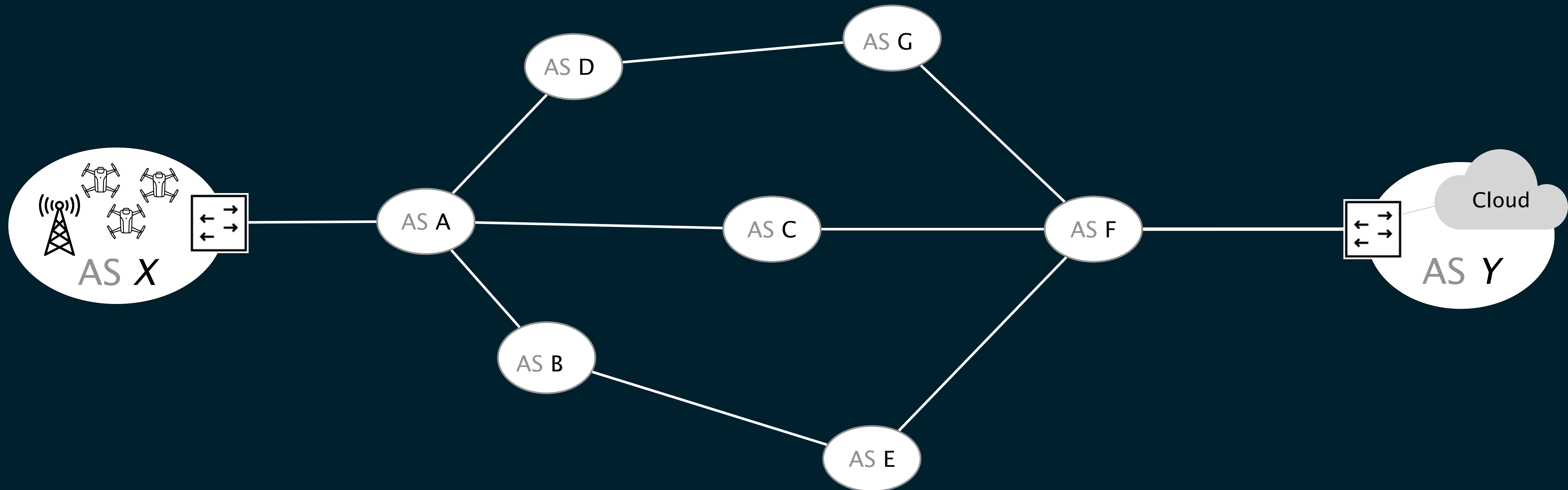
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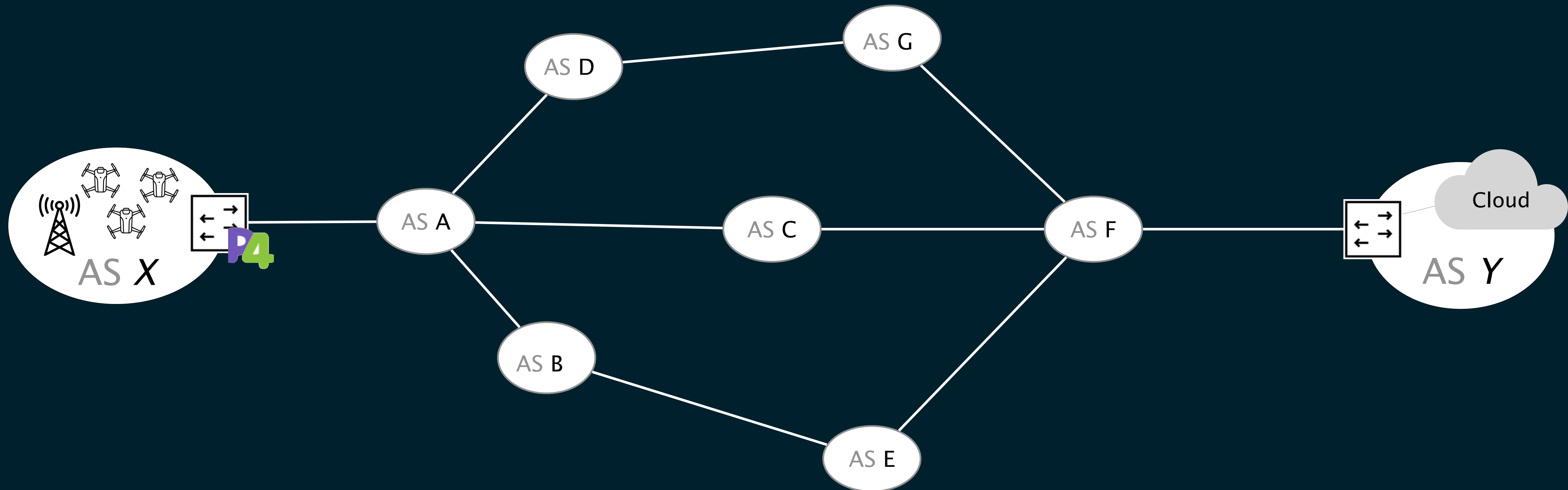


To communicate with ASY, ASX



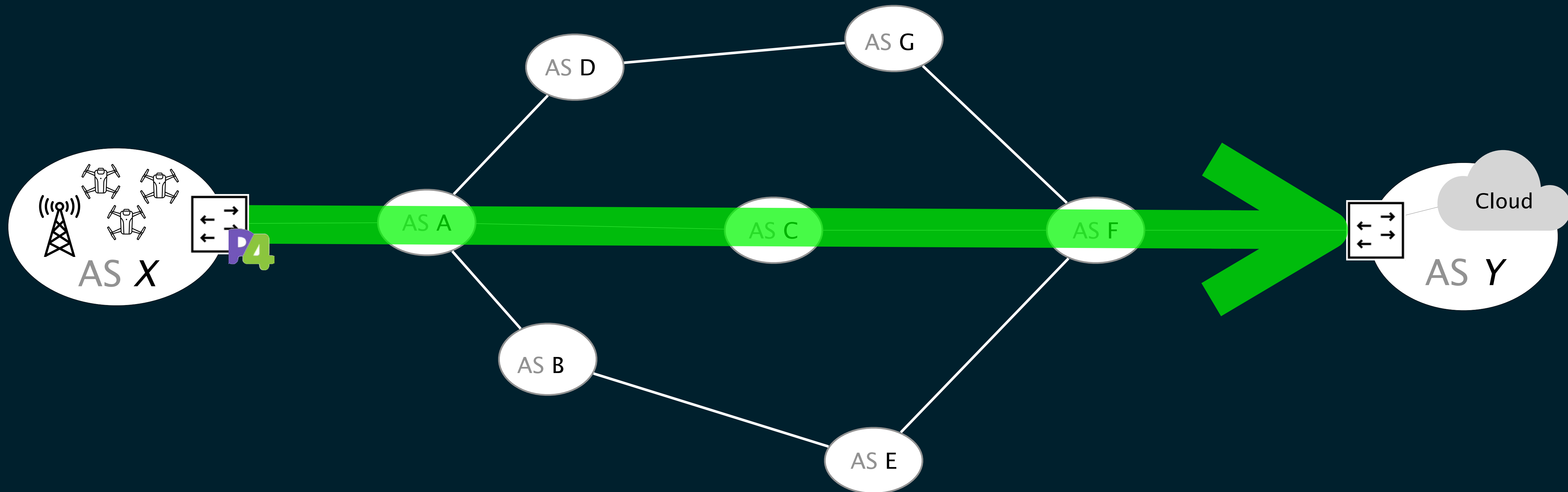


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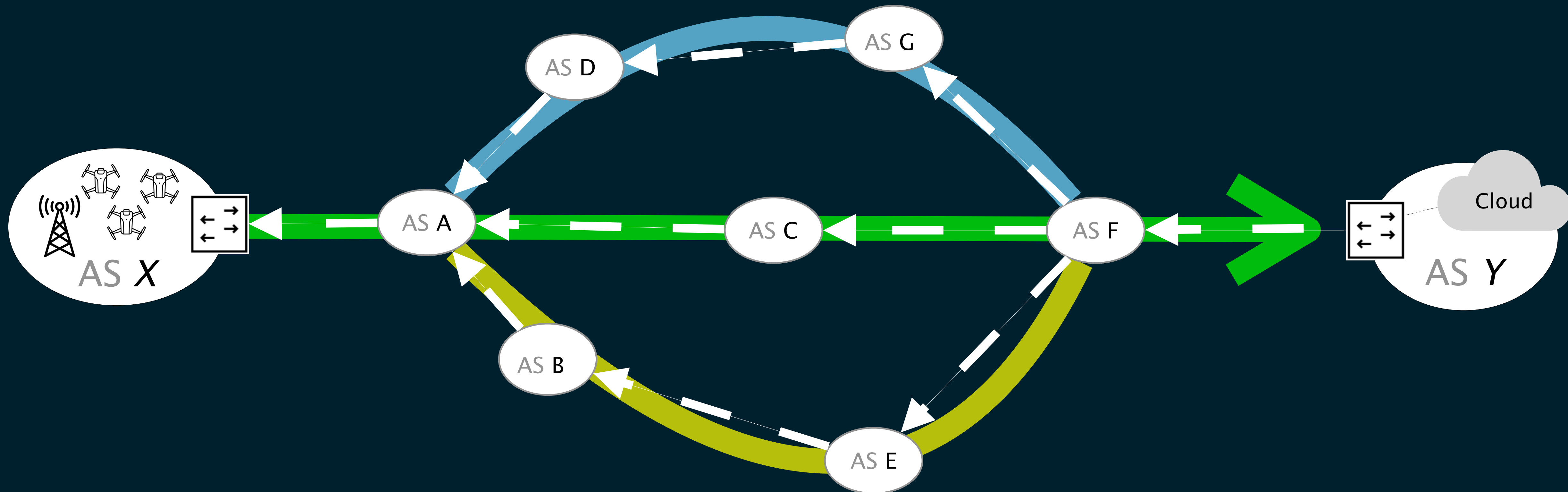




To communicate with ASY, ASX can only use one path



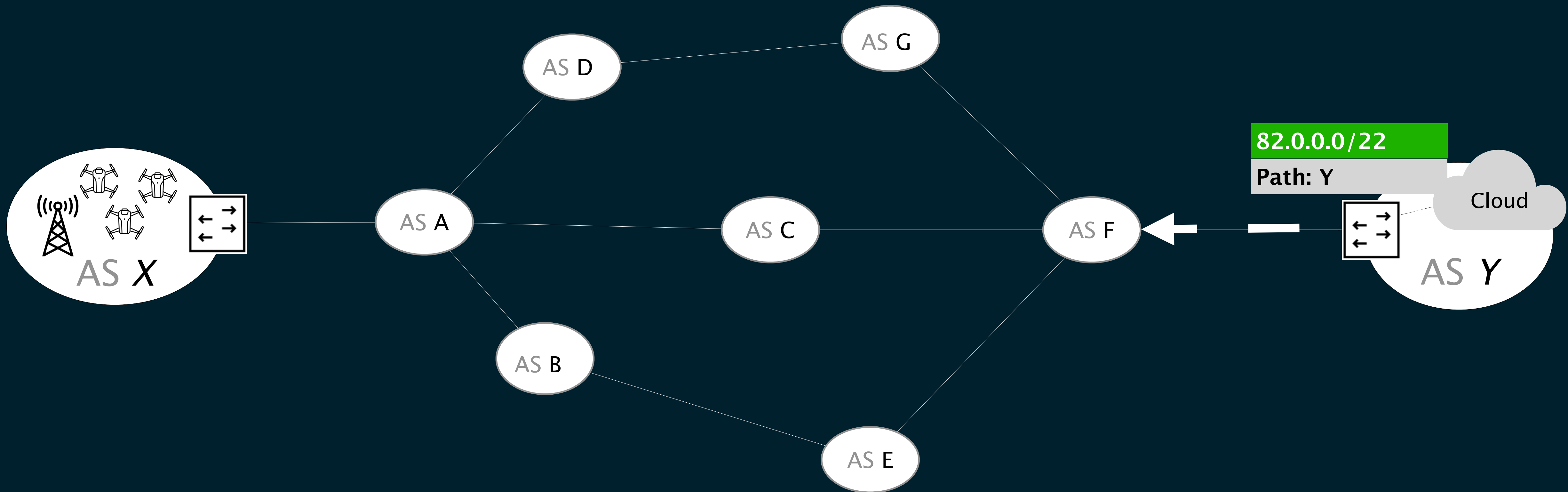
To communicate with ASY, ASX can only use one path despite the path diversity, and independently of performance





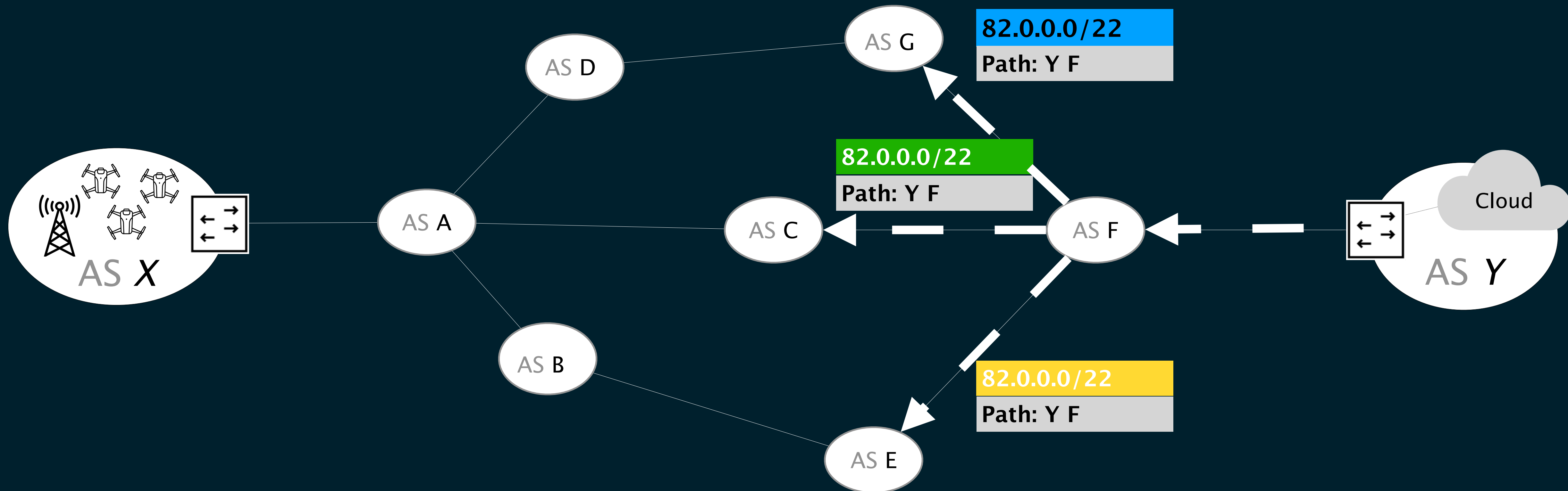




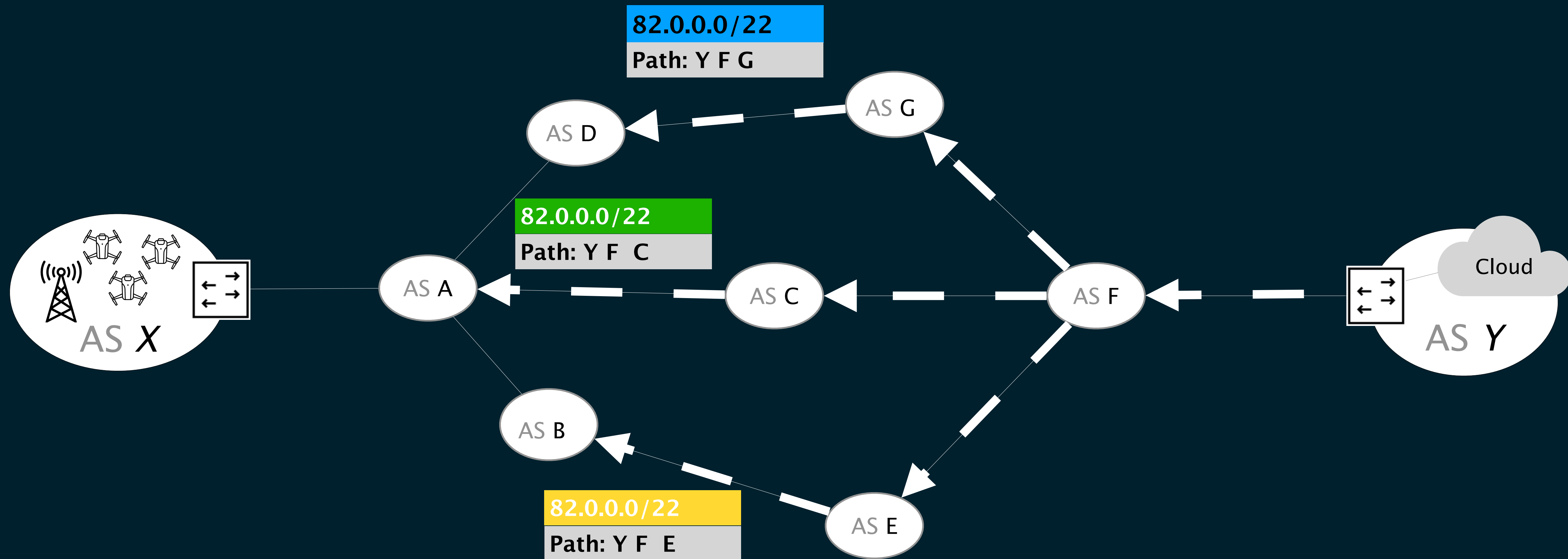




# The BGP advertisement is propagated via multiple paths

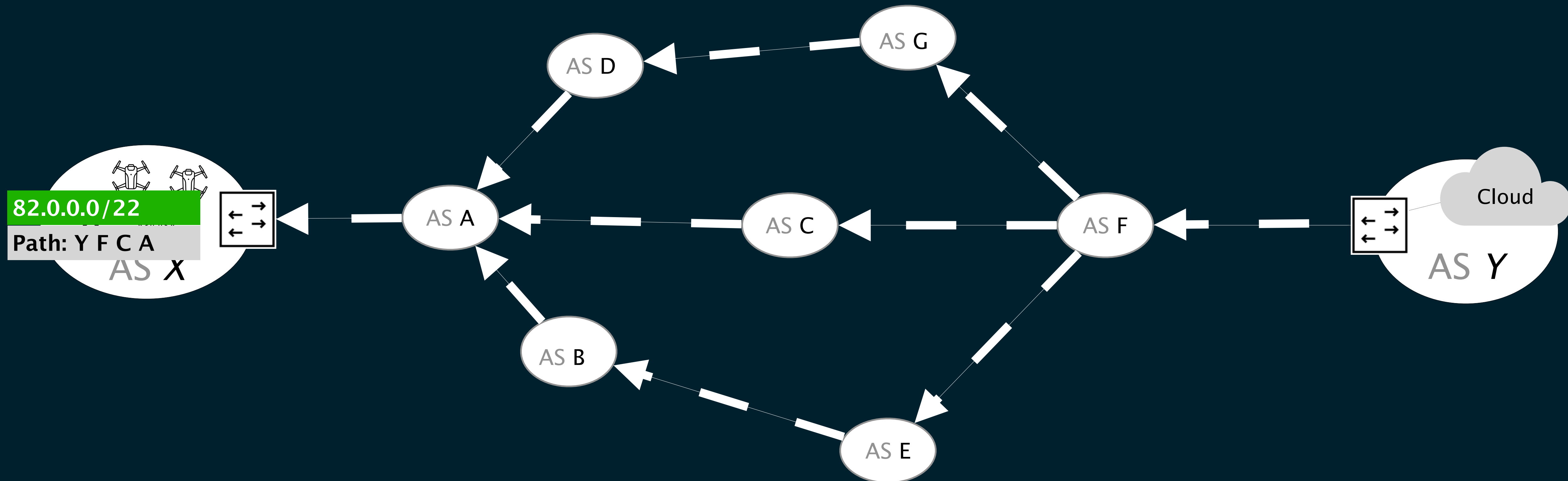


# The BGP advertisement is propagated via multiple paths





The BGP advertisement is propagated via multiple paths  
But only a single advertisement reaches the sender



What can you do with a couple of programmable points in the Internet?

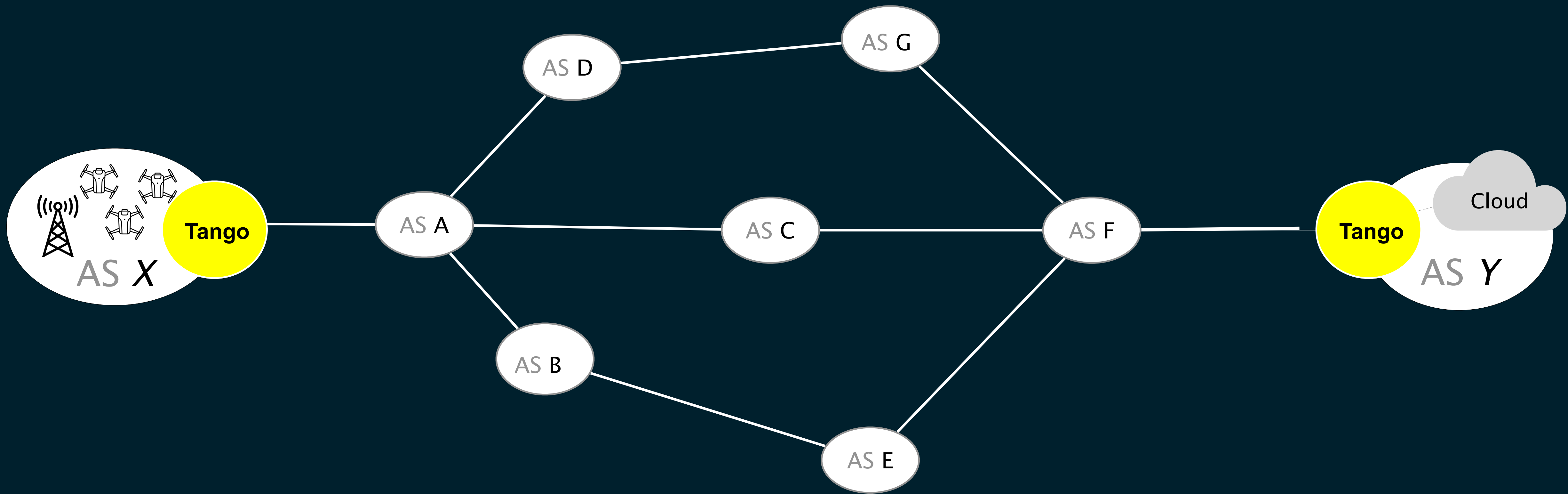
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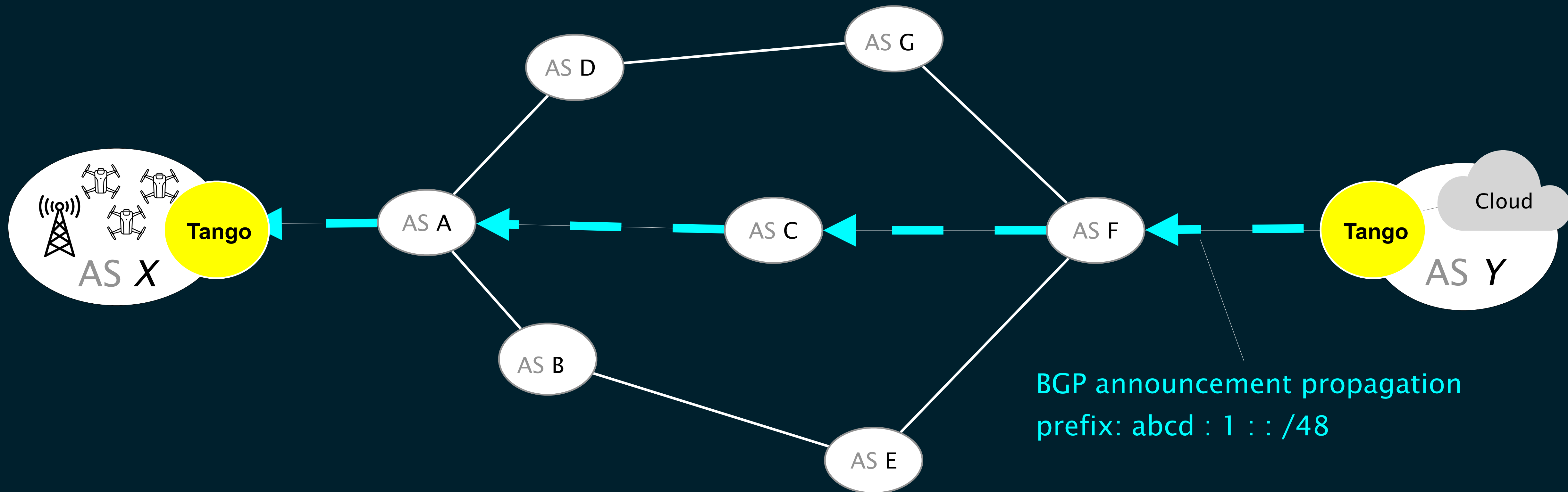
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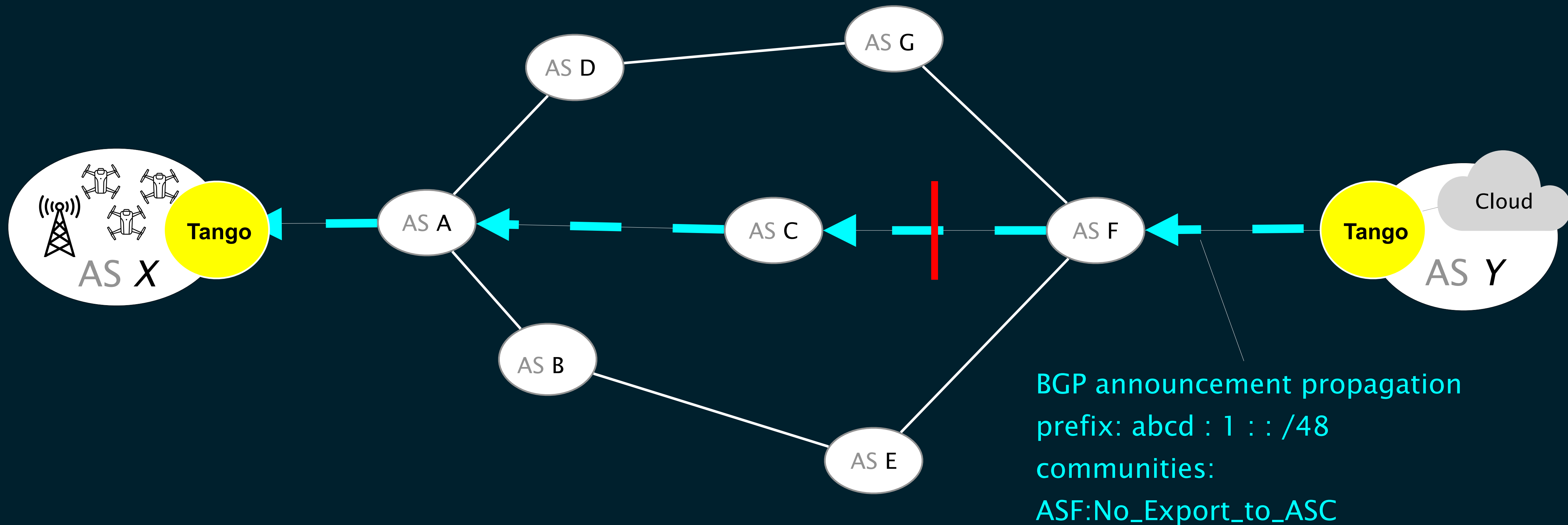


ASX only sees a single path, exported by its upstream AS

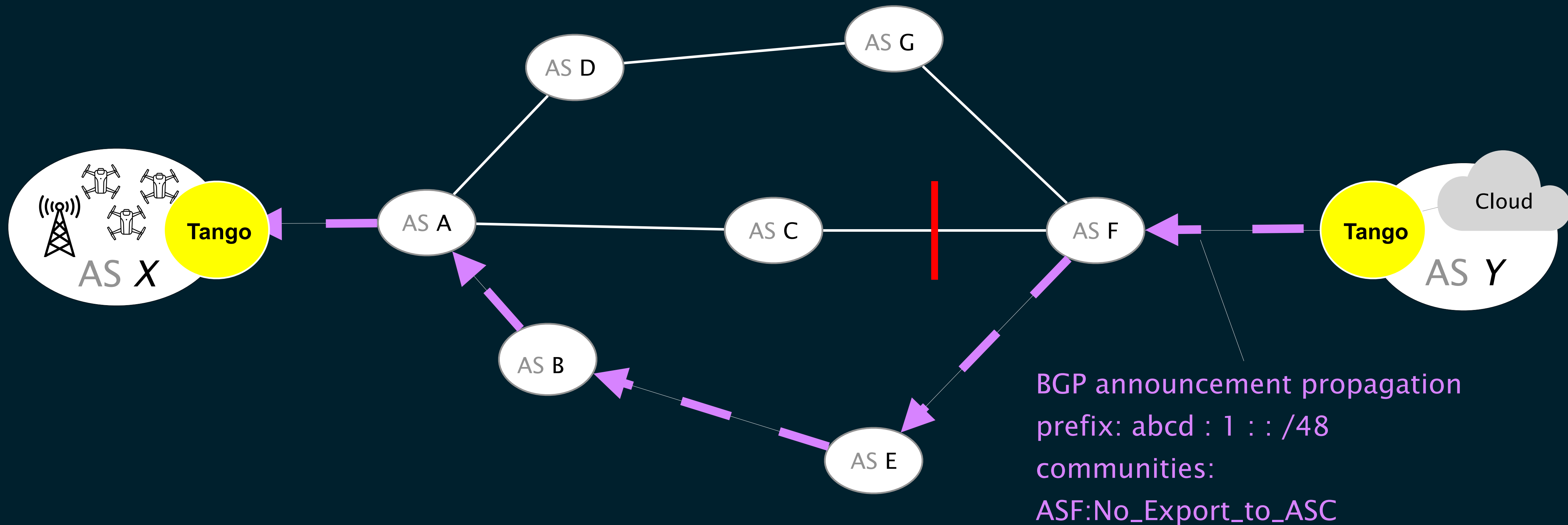




The Tango receiver advertises its IP prefix  
while suppressing the propagation of the default path

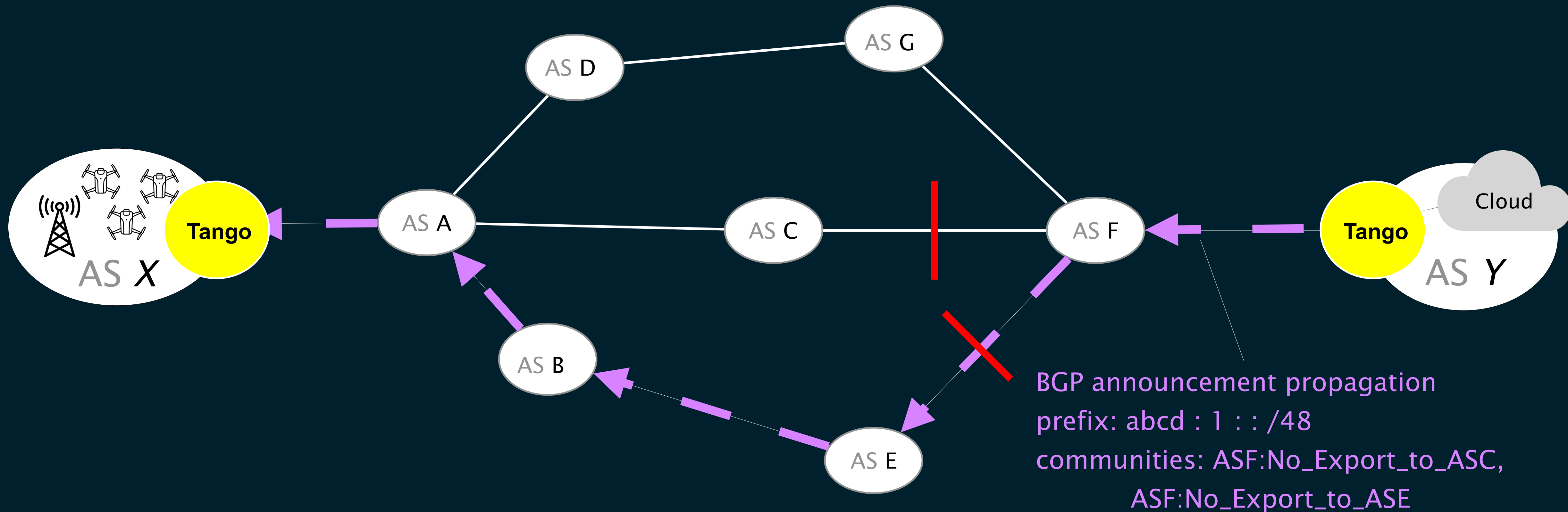


## The Tango receiver finds a new path through AS E

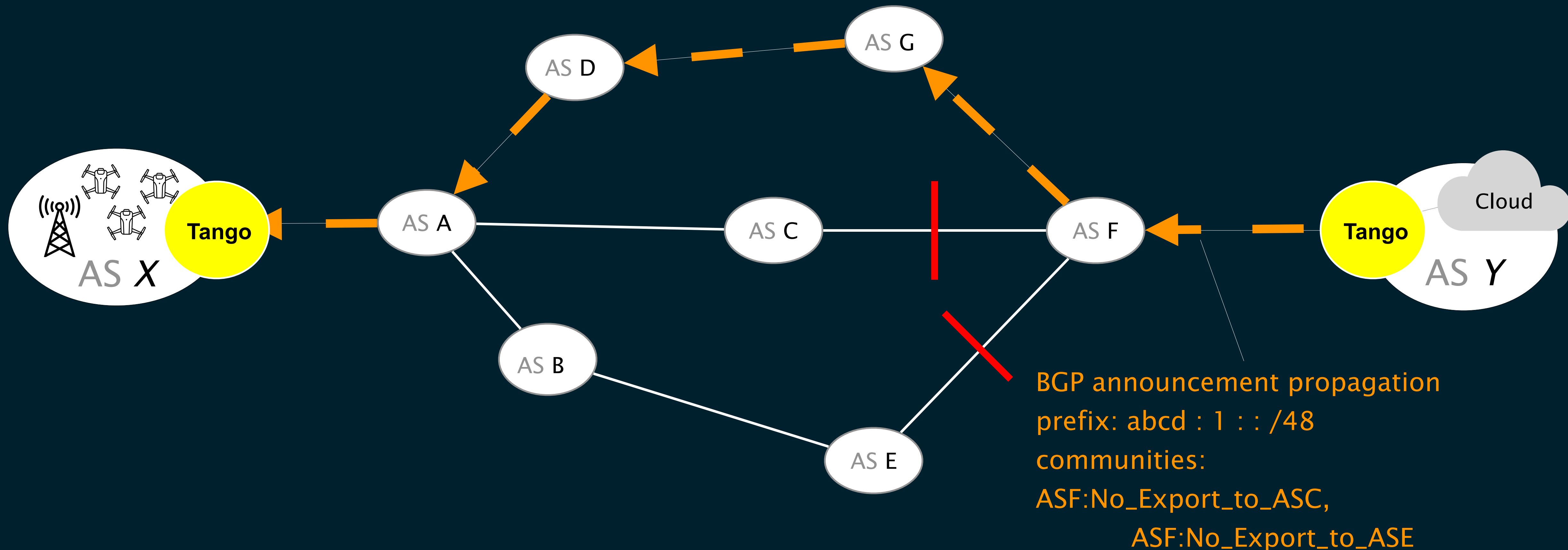




The Tango receiver finds a new path through AS E which it will again suppress

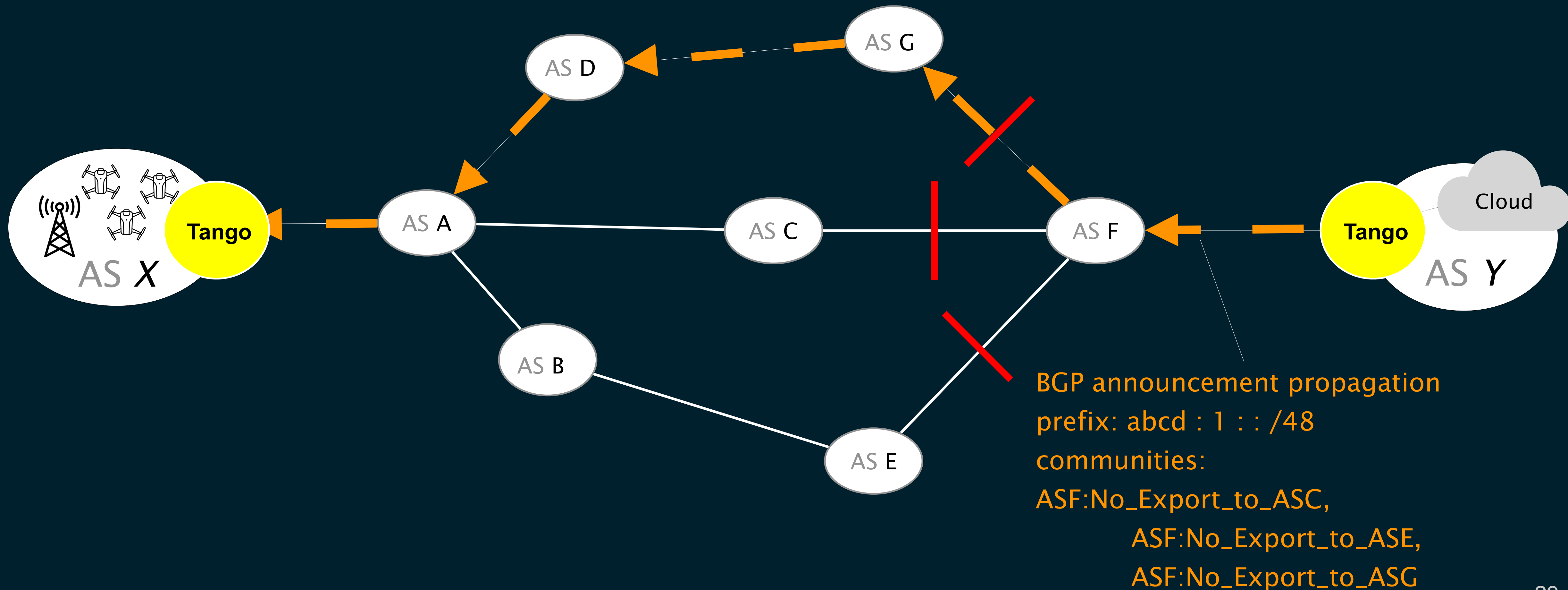


The Tango receiver finds a new path through AS E  
which it will again suppress to find yet another path through AS G

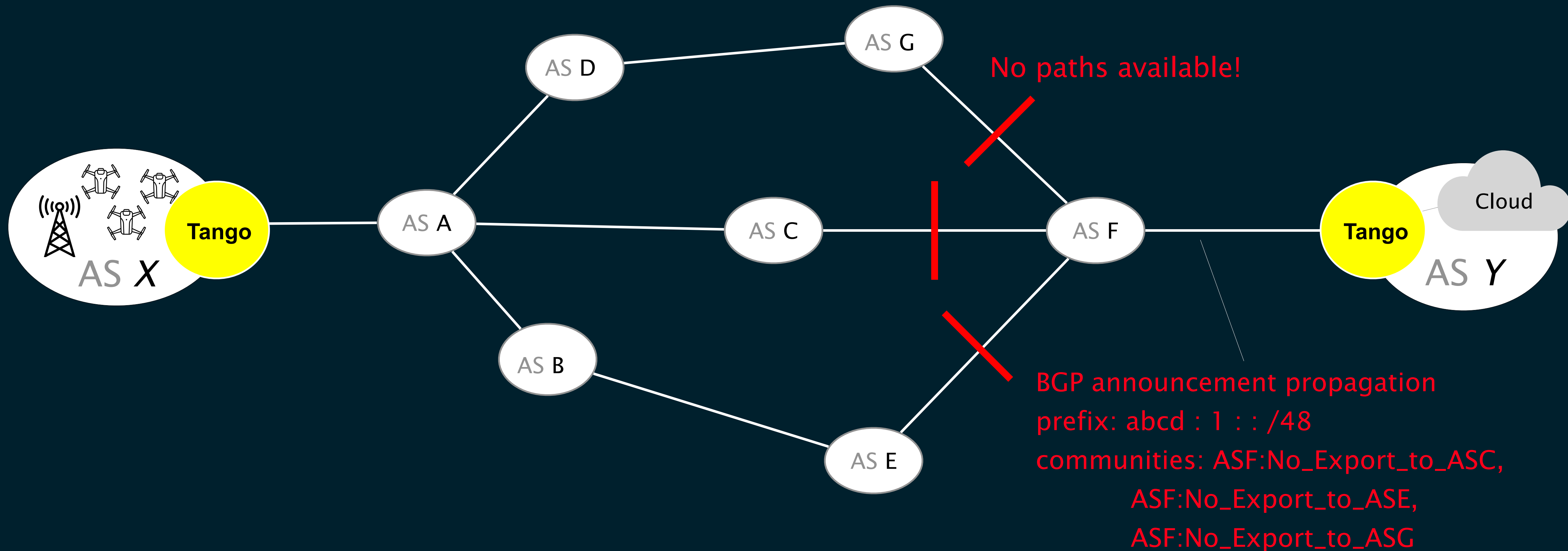




The Tango receiver stops when there are no new paths

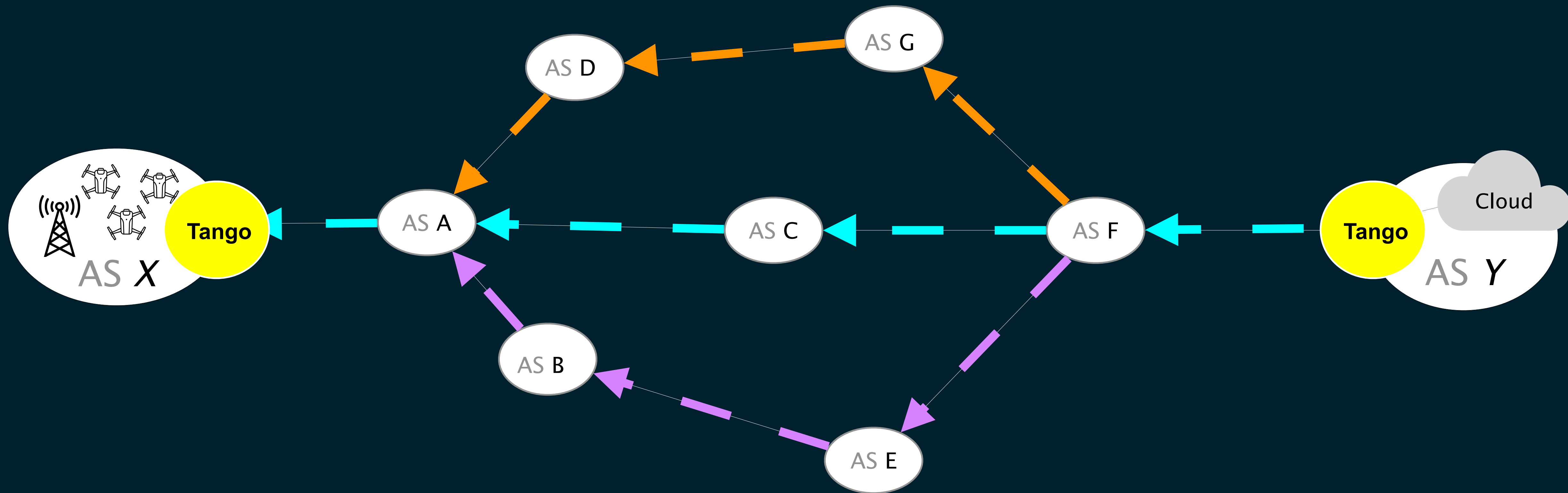


# The Tango receiver stops when there are no new paths





## AS Y announces different IP prefixes along different paths









# Global testbed

Run Tango– Pathfinder from 23 nodes hosted by Vultr to exposes Internet paths

Routed traffic over the exposed and default paths to two destinations: LA and Stockholm

Collected latency and loss measurements every 10ms, over roughly 32 hours



# Tango-paths outperform the default path

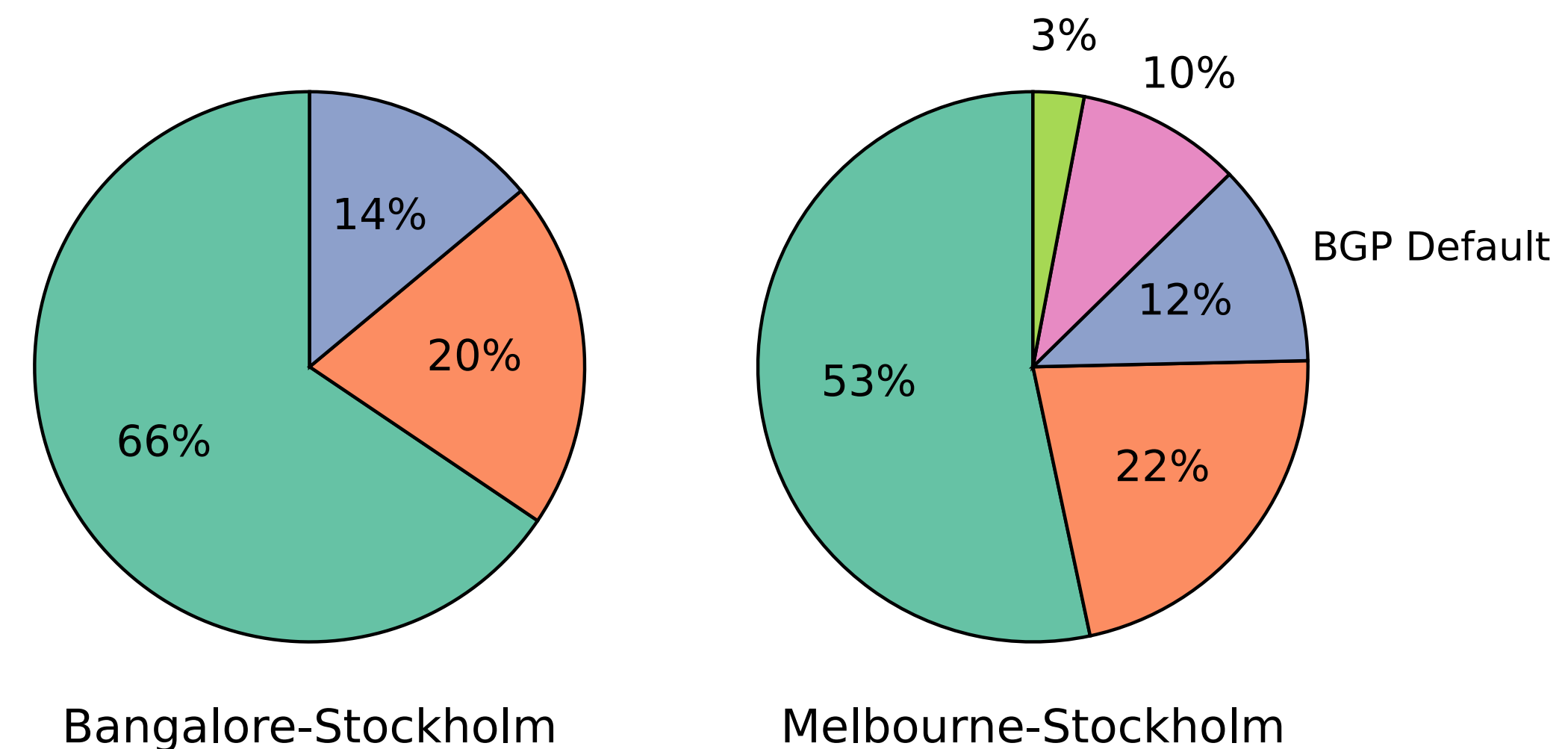
Across 23 measured pairs, 20 pairs had alternative paths that outperformed the default:

100% of the time for 15 pairs  
75–88% of the time for 5 pairs

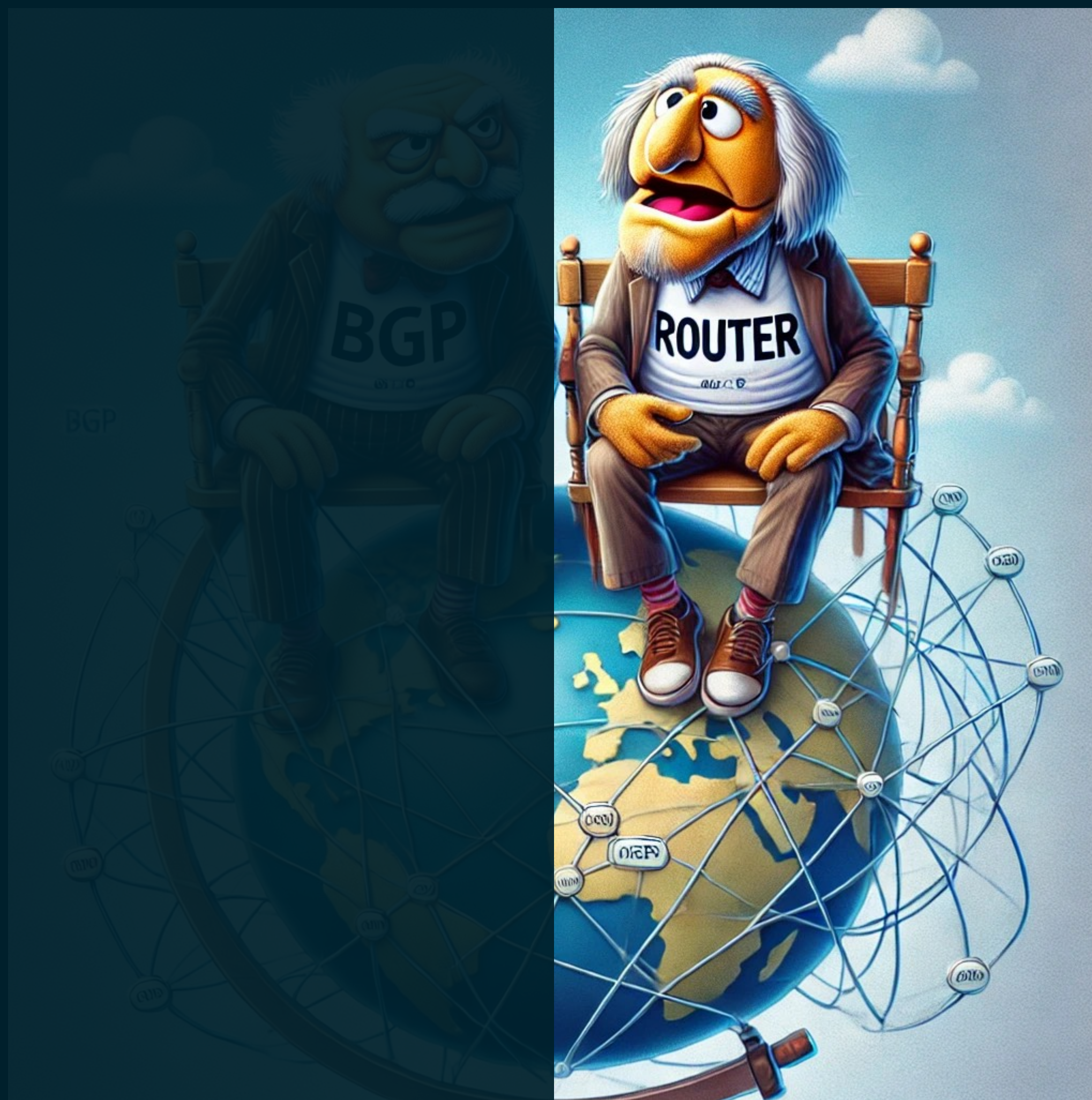
Bangalore–Stockholm: BGP default beaten by alternative paths **100% of the time**

Melbourne–Stockholm: BGP default beaten by alternative paths **88% of the time**

Breakdown of best paths for two pairs









# Tango's design requirements for performance-driven routing

## Route Control

Tango senders need to control which path traffic will use.

## Accurate Measurements

Measurements should not be affected by irrelevant conditions e.g., slow receivers, Wi-Fi.

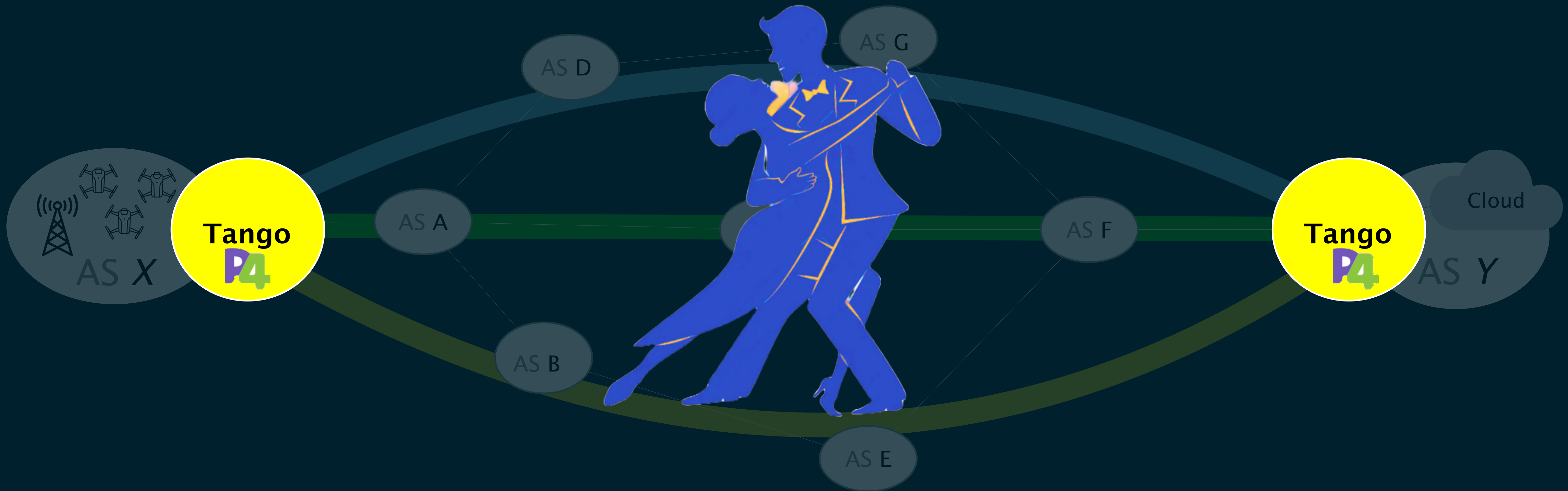
## Trustworthy Measurements

An on-path attacker should not be able to distort measurements to their advantage.

## Dynamic & Secure Rerouting

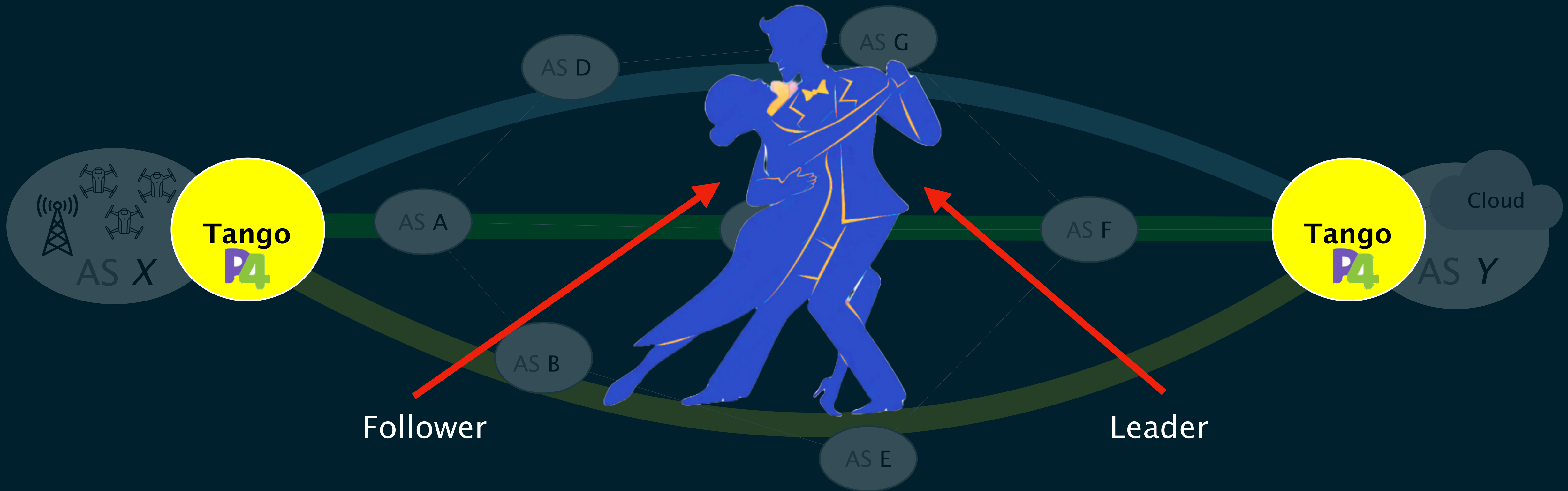
Tango should allow dynamic performance-driven and safe reroutes.

Tango solves these challenges with P4 and co-operation

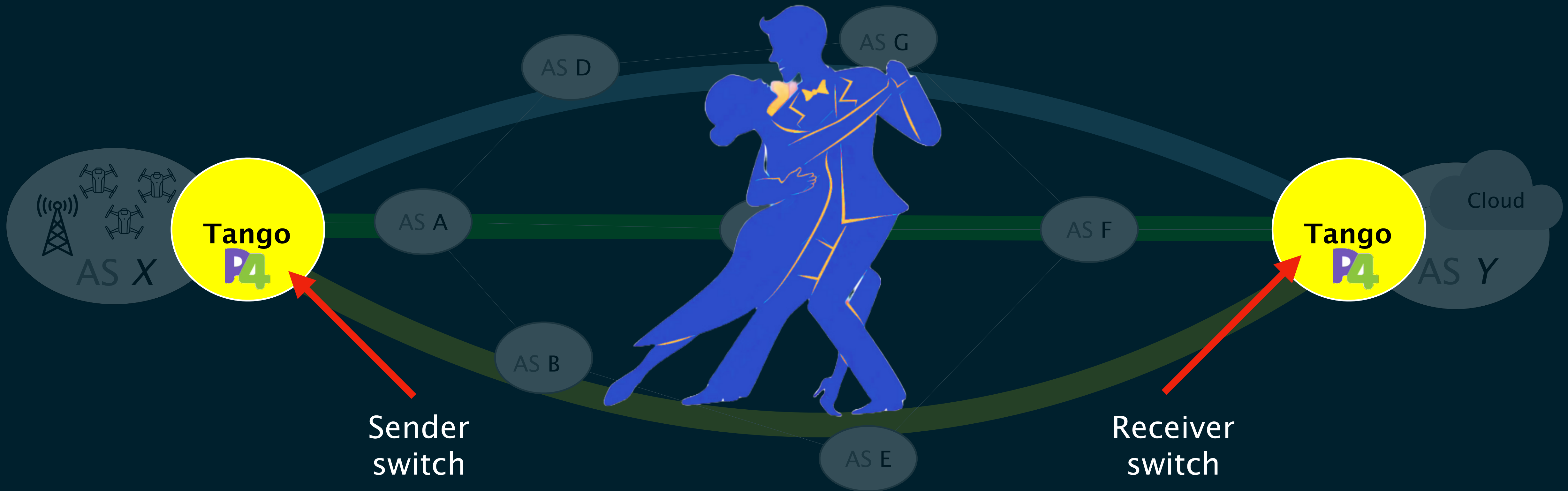




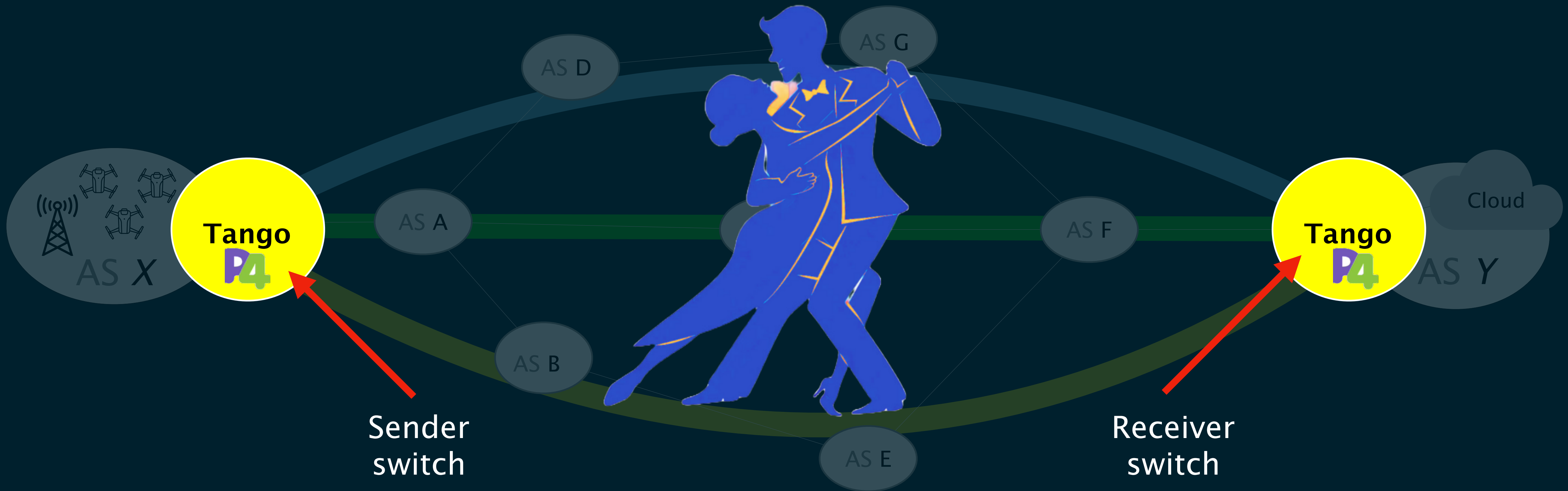
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The sender switch performs the move that the receiver has signaled.



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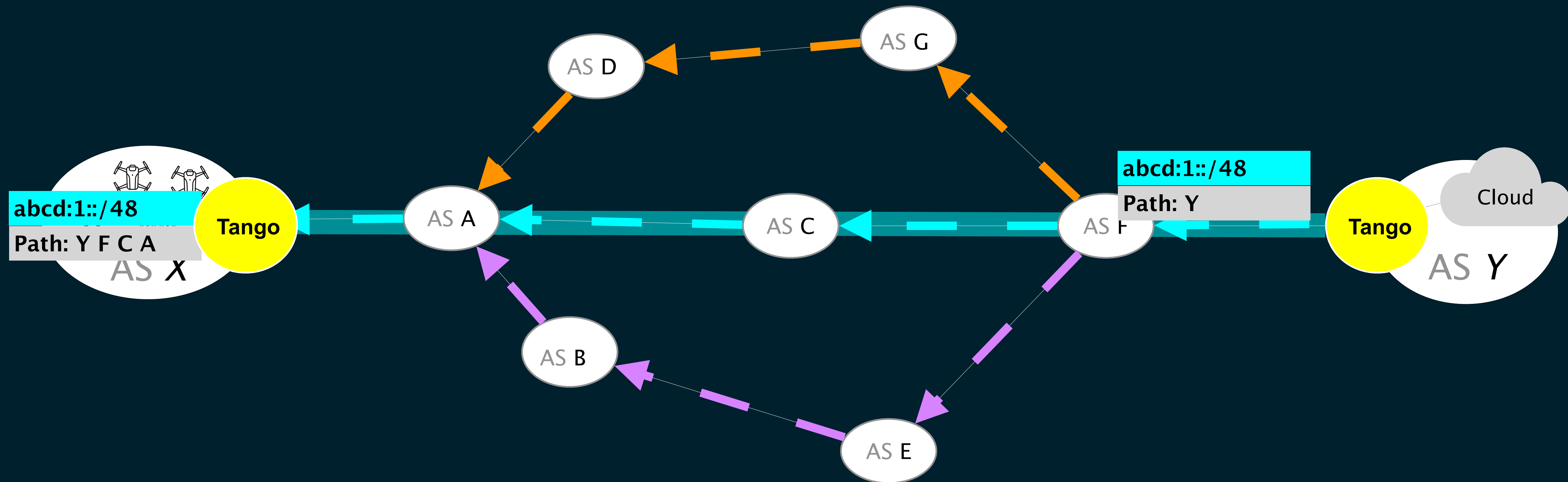
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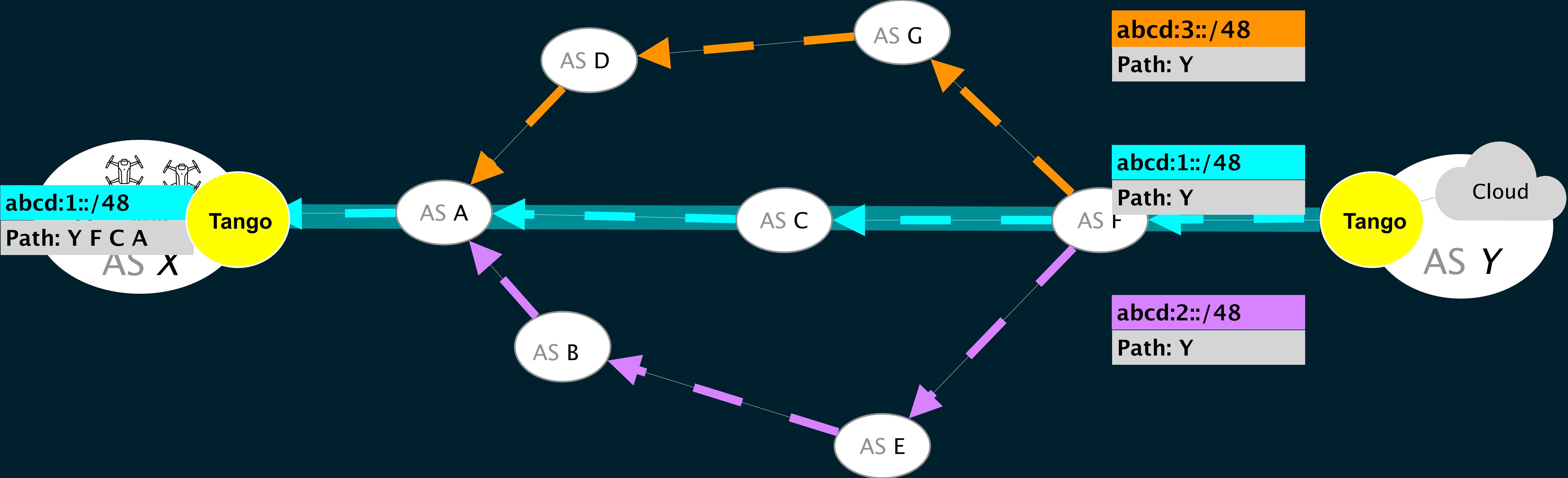
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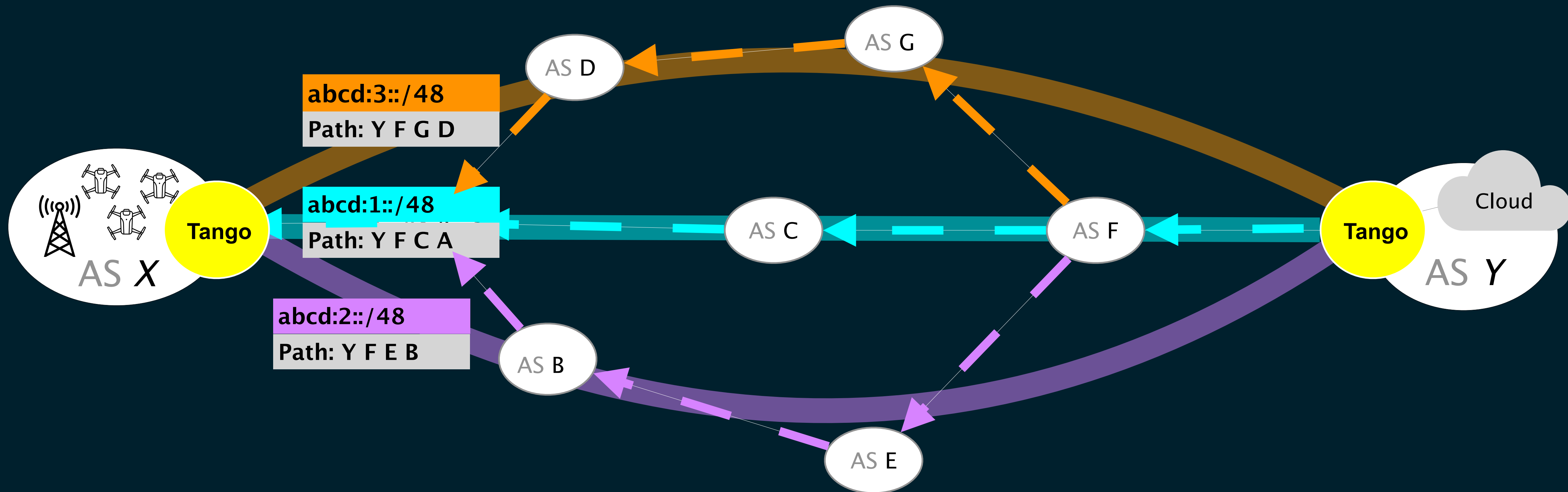


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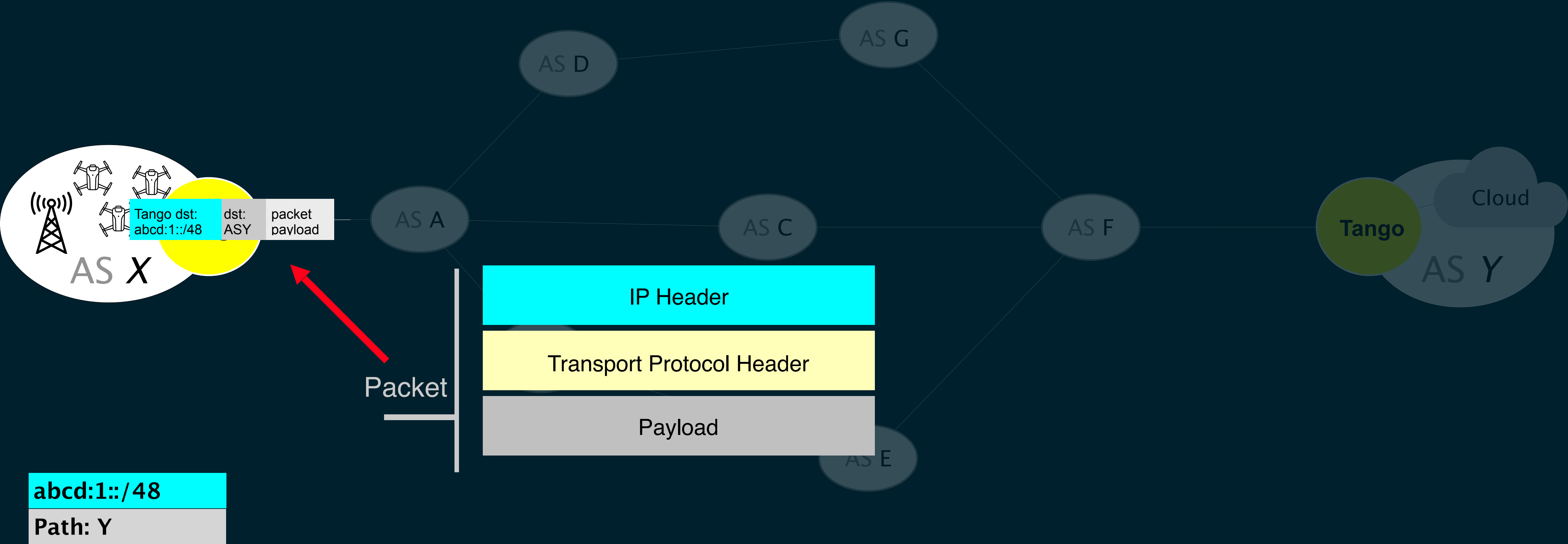




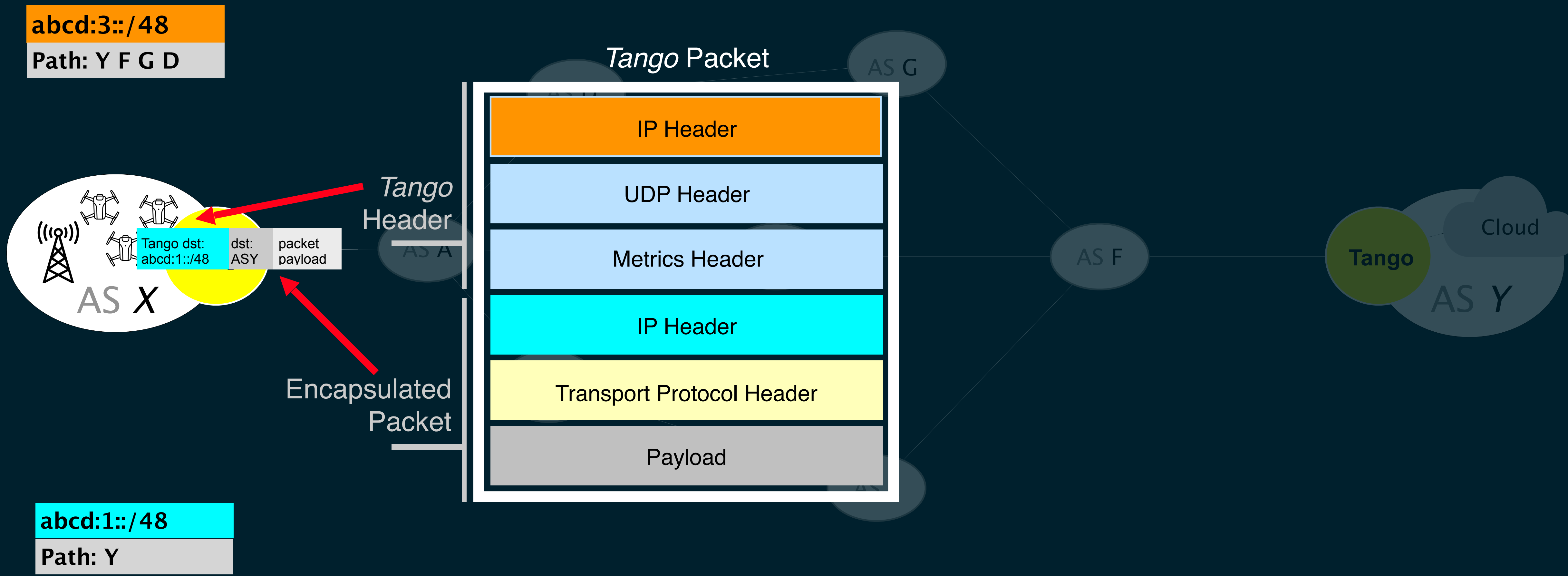
## AS Y announces different IP prefixes along different paths



Upon reception of a packet,

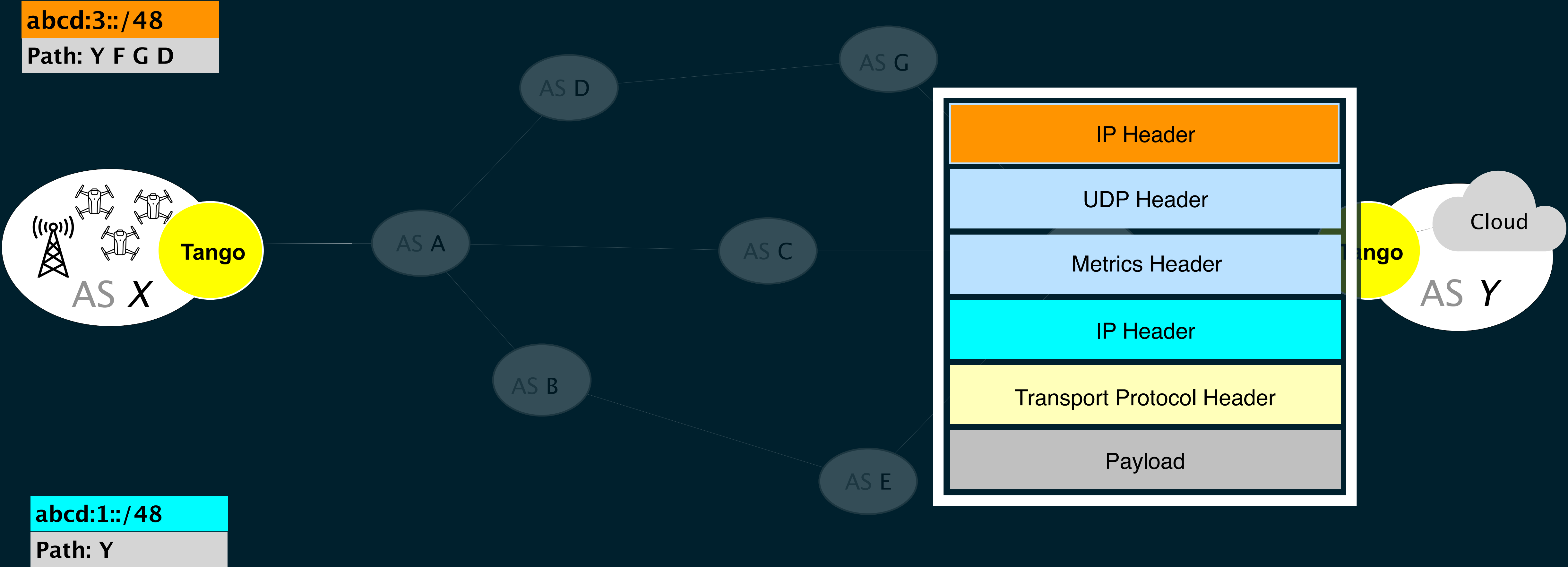


Upon reception of a packet, the sender encapsulates it with a destination within the prefix that correspond to the path of choice

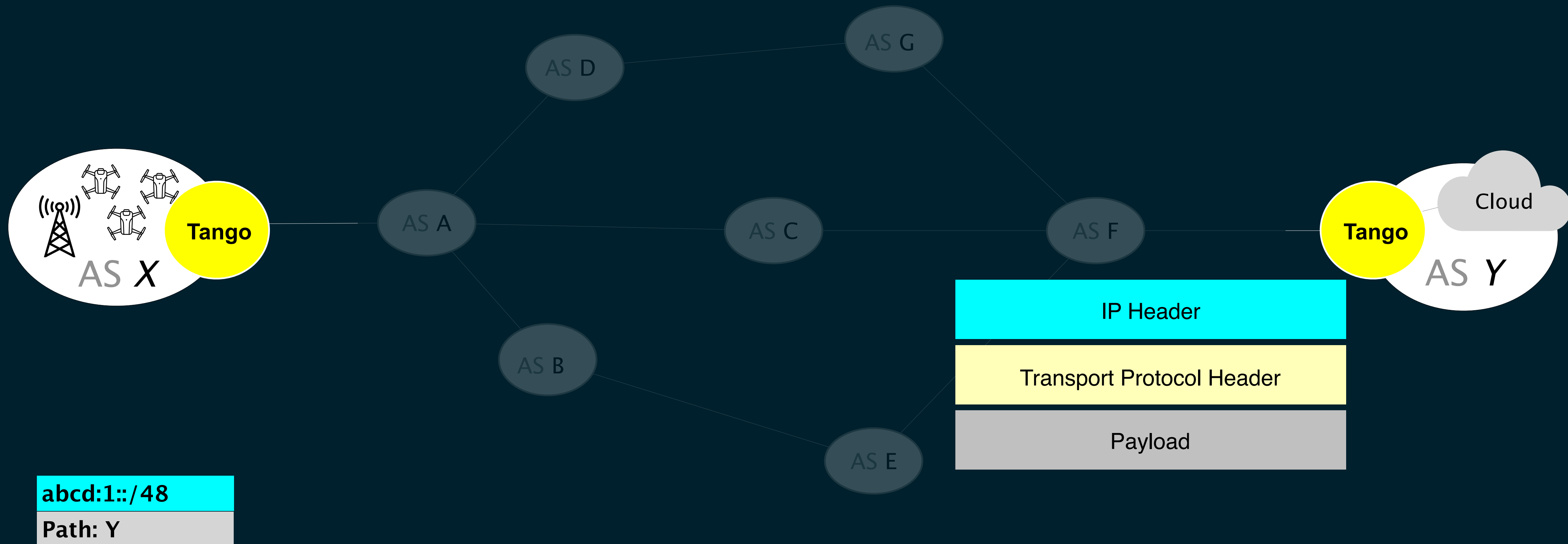




The receiver decapsulates packets before letting them reach their destination



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# Tango's design requirements for performance-driven routing

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## Trustworthy Measurements

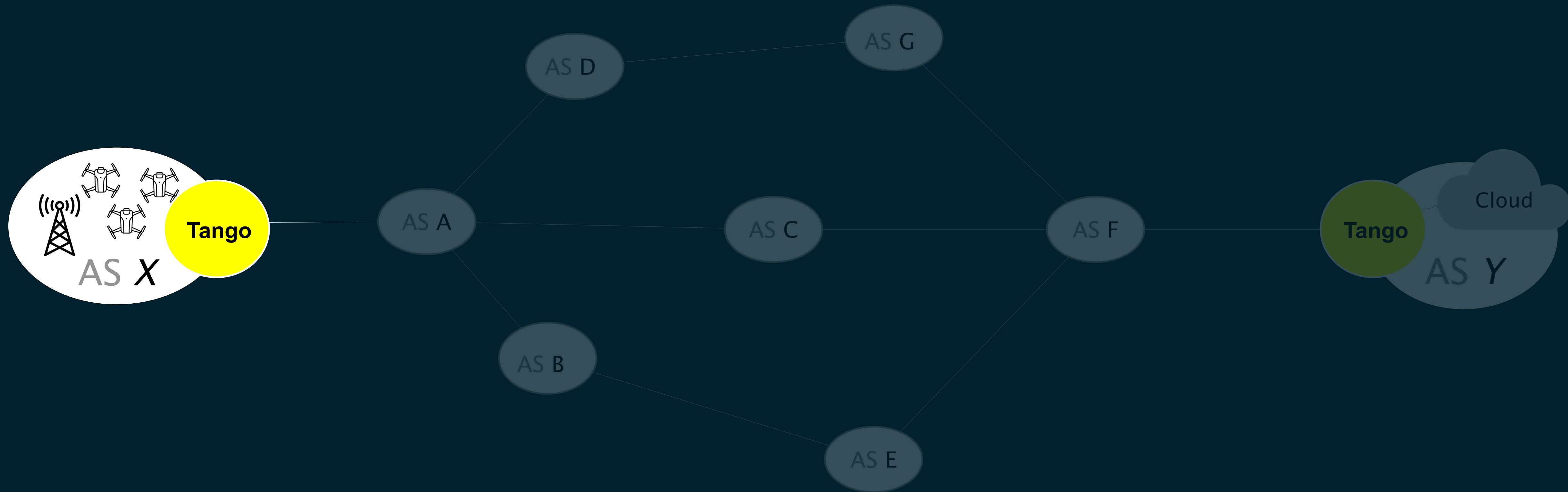
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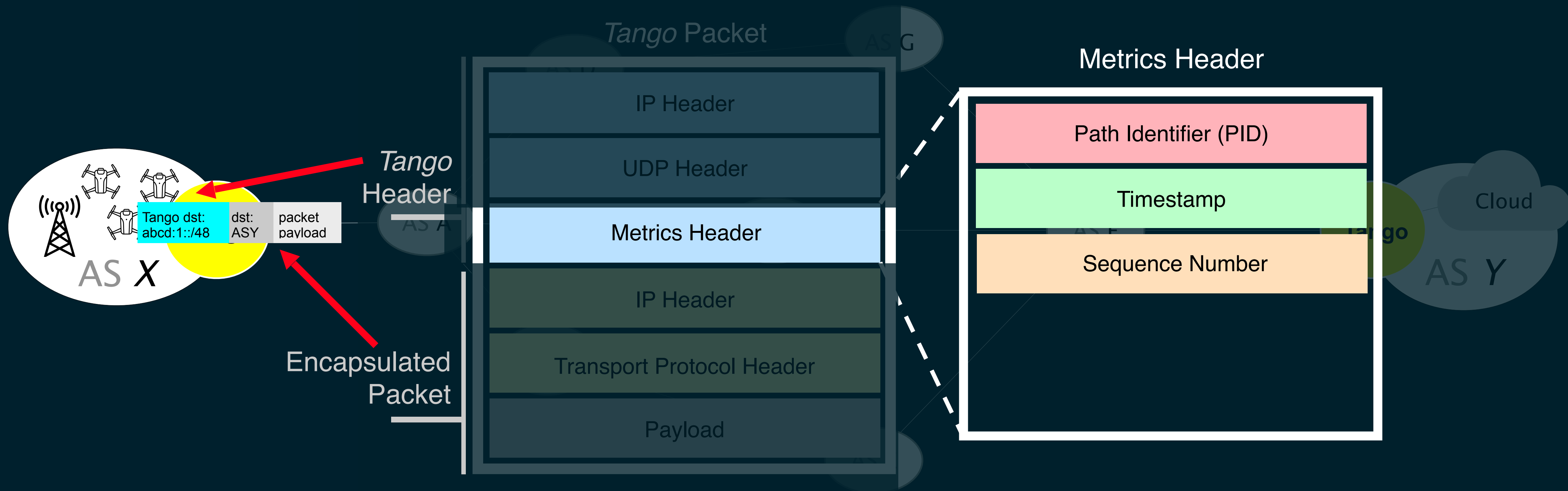
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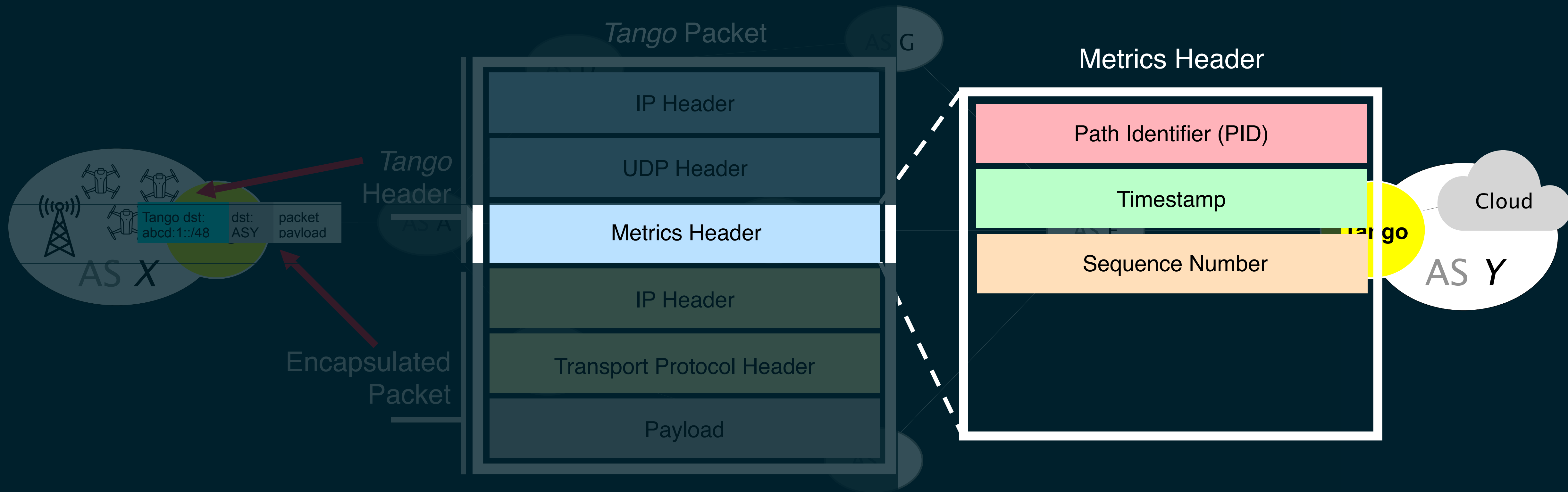
Conventional active round-trip measurements are inaccurate and can be easily manipulated



The sender includes the timestamp of each packet's departure, and a per-path sequence number in the Tango header.



The receiver calculates one-way latency and loss for each path avoiding the noise of the access networks





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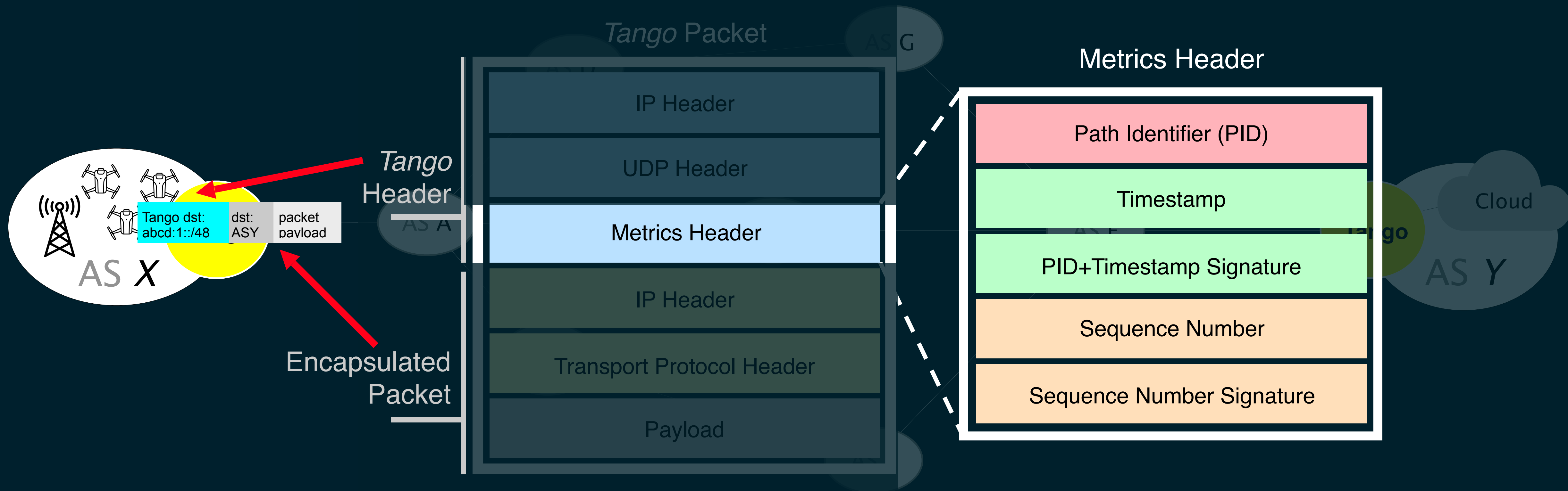
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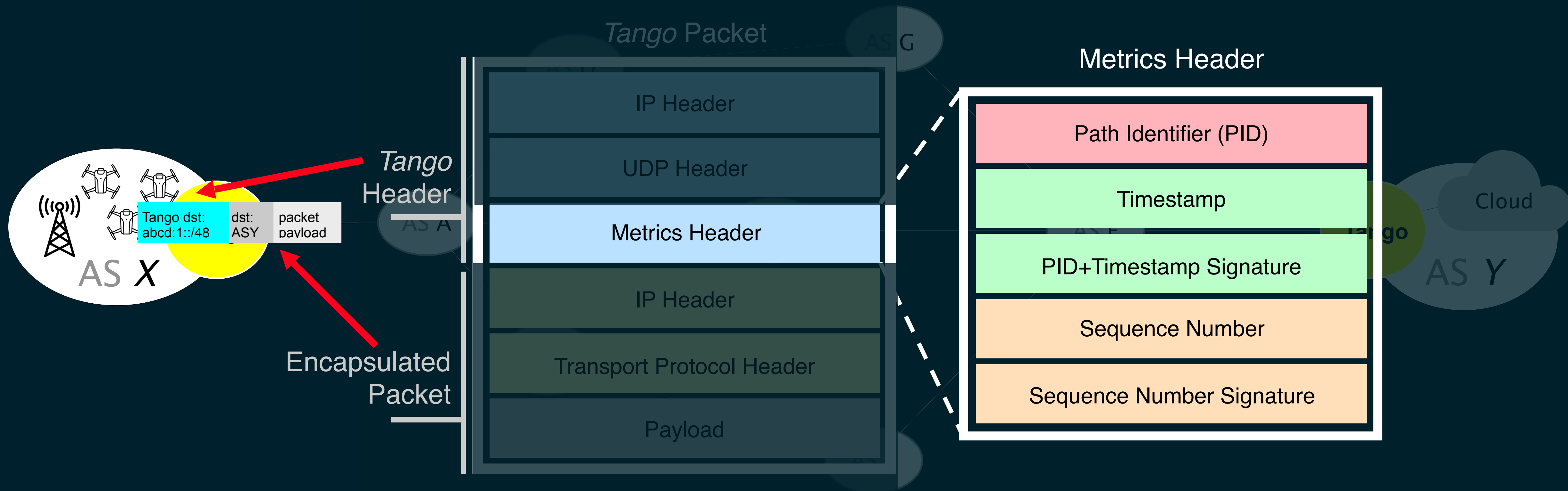
## Dynamic & Secure Rerouting

Tango should allow dynamic performance-driven and safe reroutes.

Tango sender adds a path-specific signature to each ms timestamp,  
an attacker cannot manipulate or replay it to affect latency measurements



Tango sender adds one bit signature to each sequence number,  
an on-path attacker would need to guess multiple to affect loss rate





# Tango's design requirements for performance-driven routing

## Route Control

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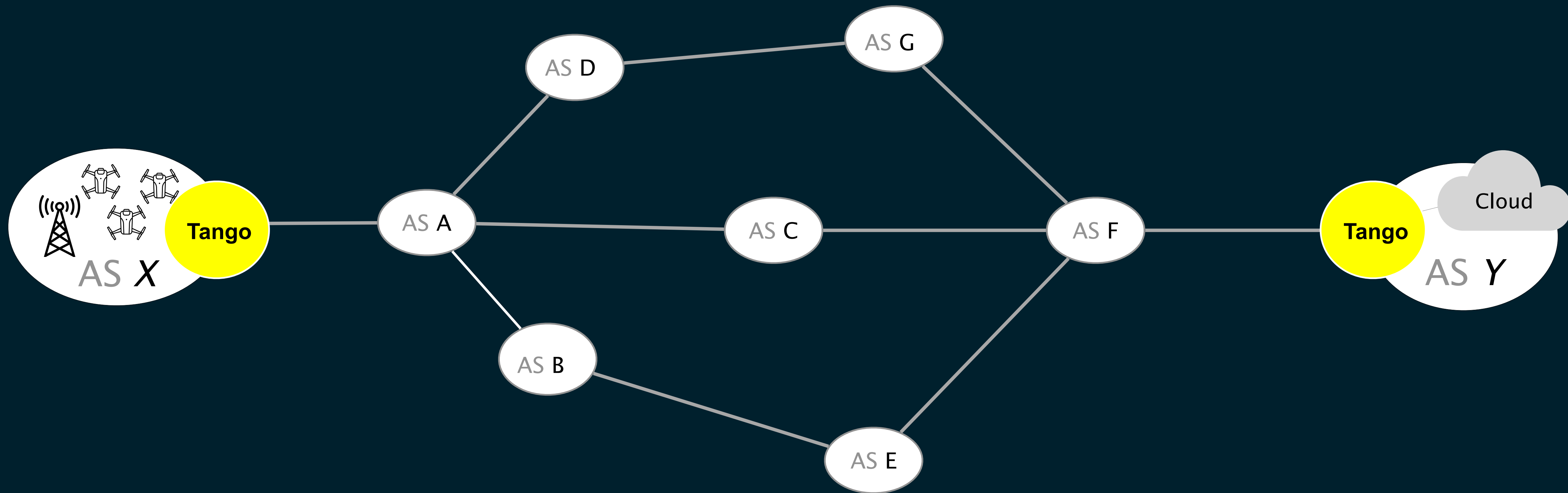
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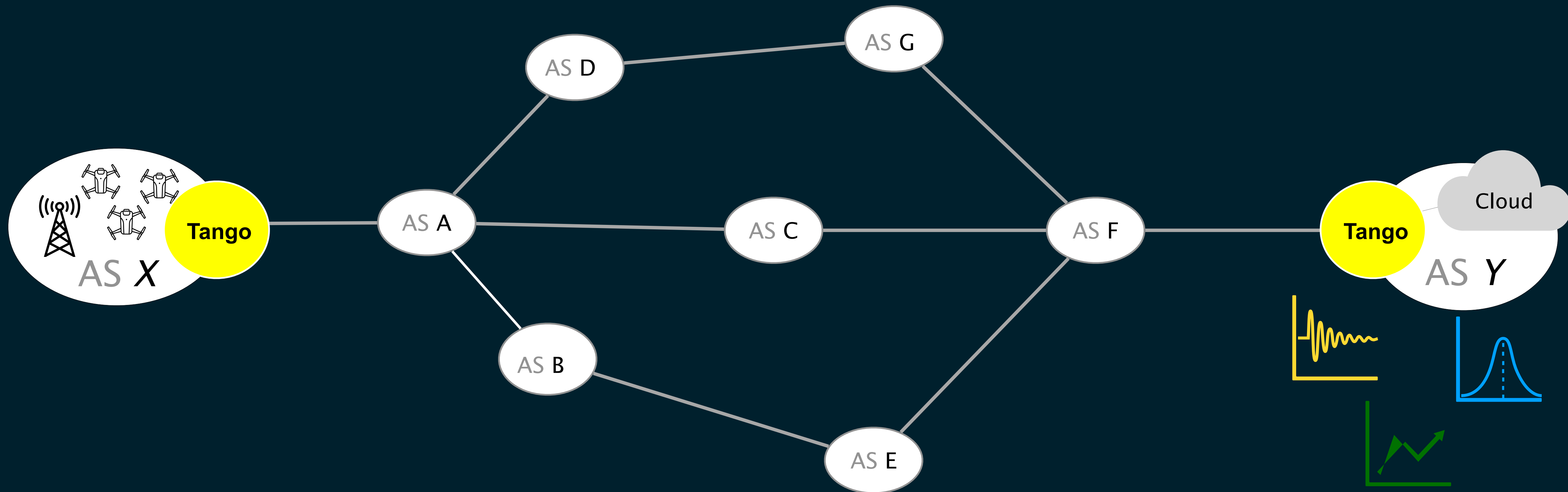
## **Dynamic & Secure Rerouting**

**Tango should allow dynamic performance-driven and safe reroutes.**

The Tango sender selects paths,

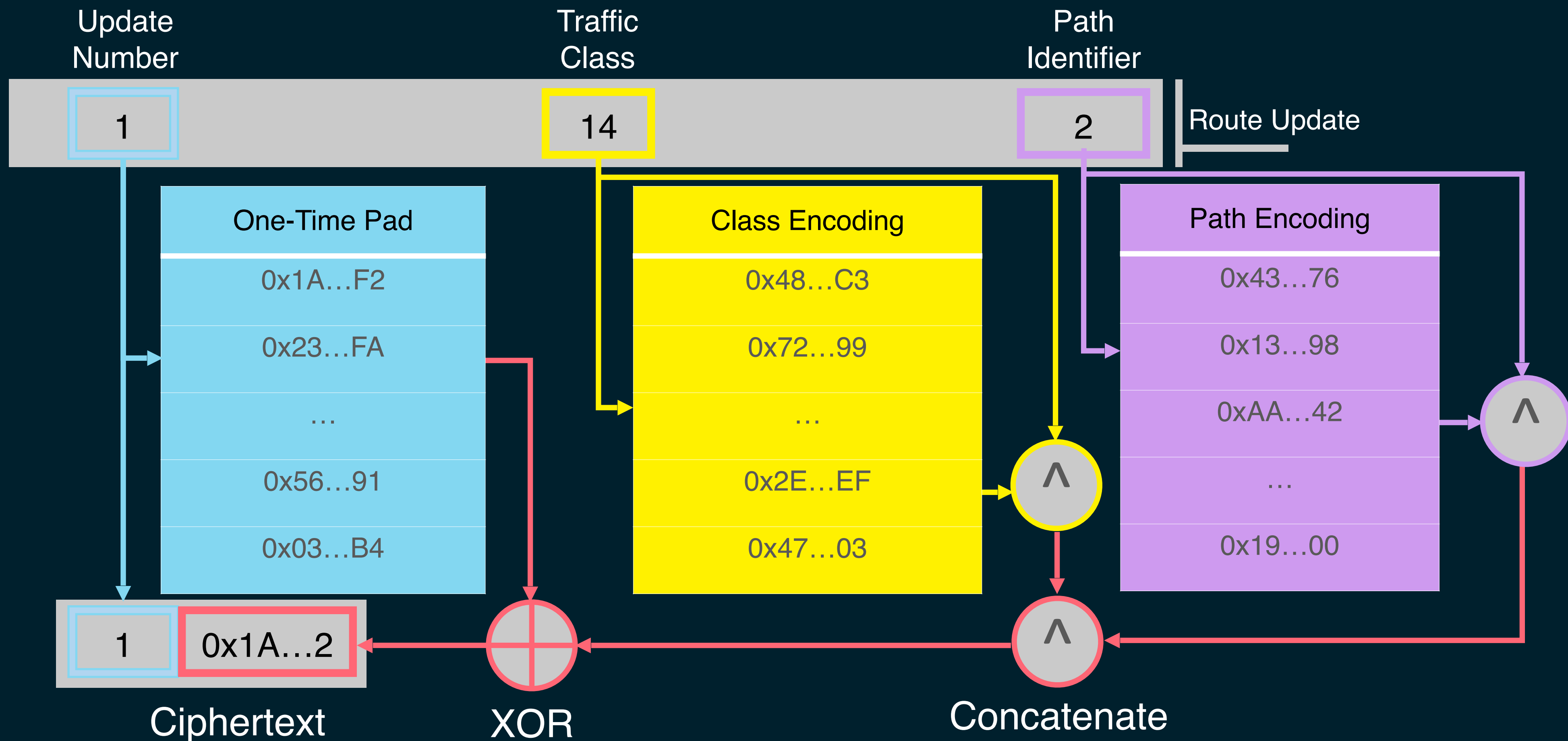


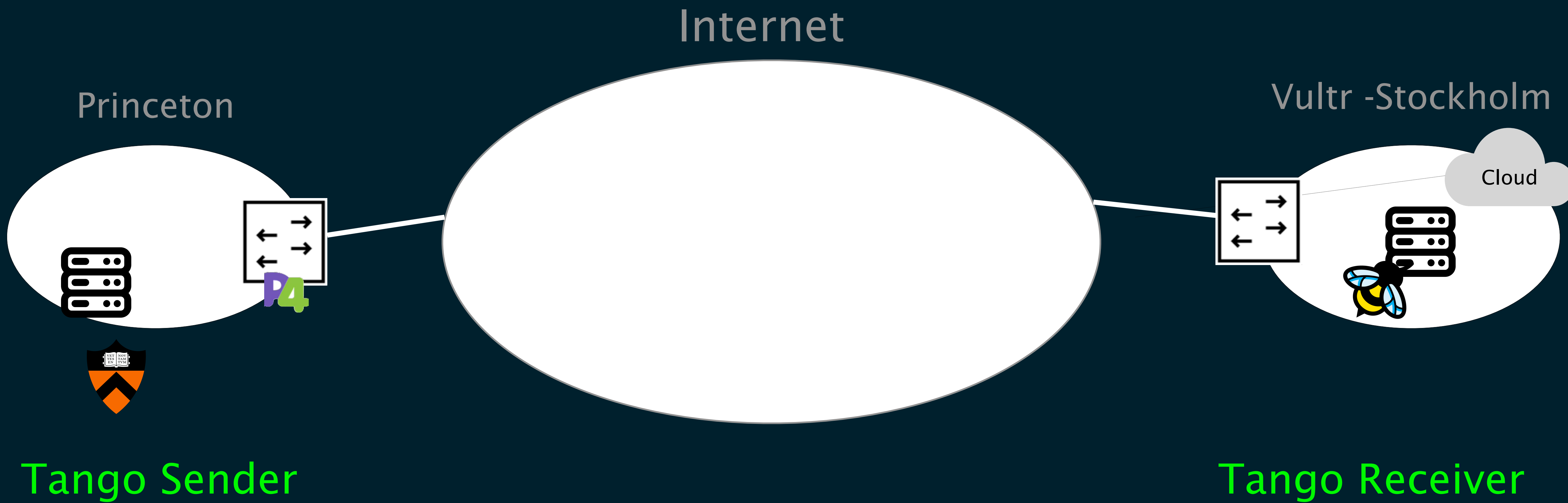
The Tango sender selects paths,  
but the Tango receiver collected the measurements





# Tango protects reroute commands with one-time-pad



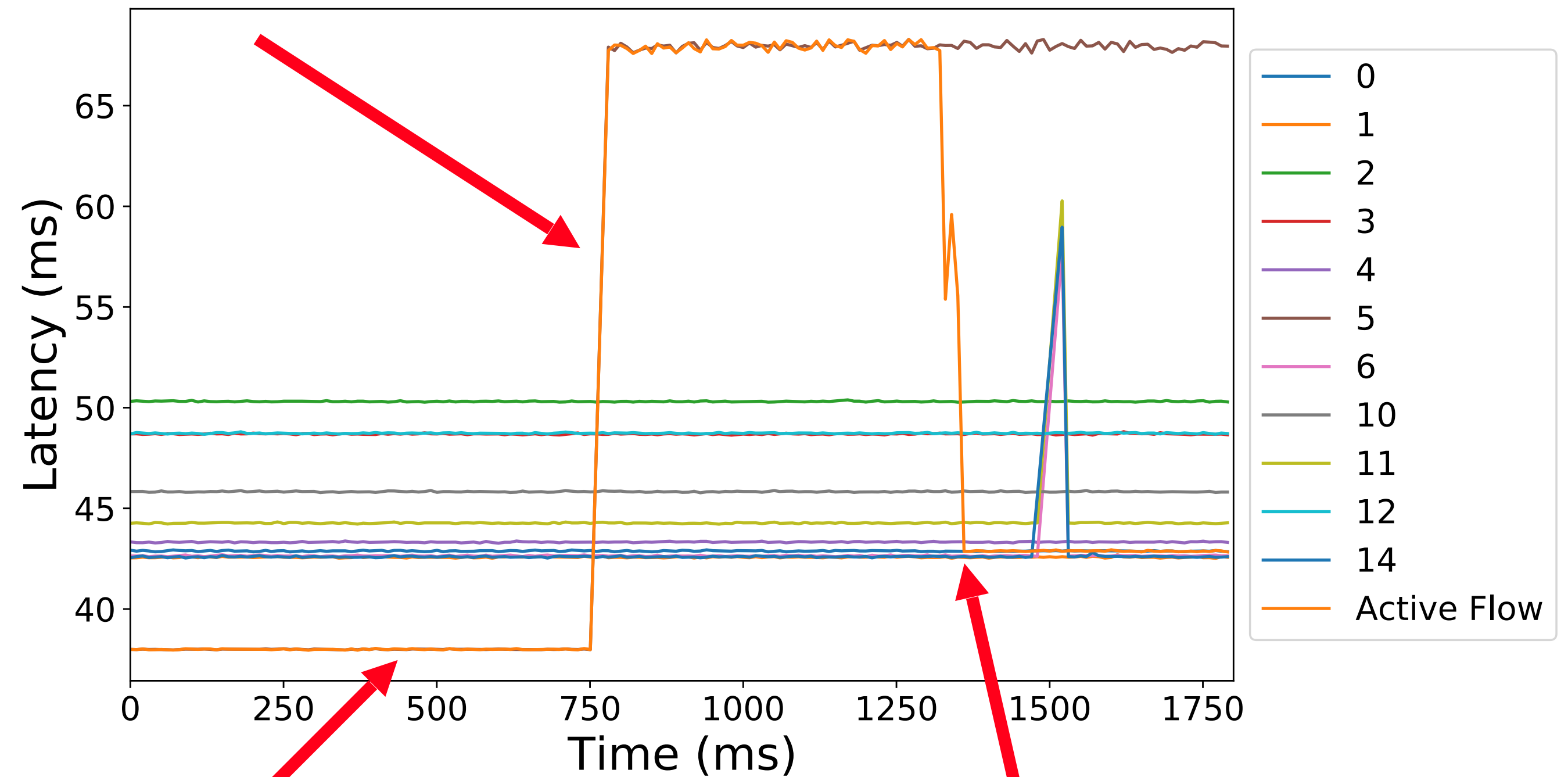


# Real-world Testbed

We run Tango between  
Princeton and Stockholm!

Route update complete in <1s

delay spike



initial best path

dynamically  
move to new  
best path

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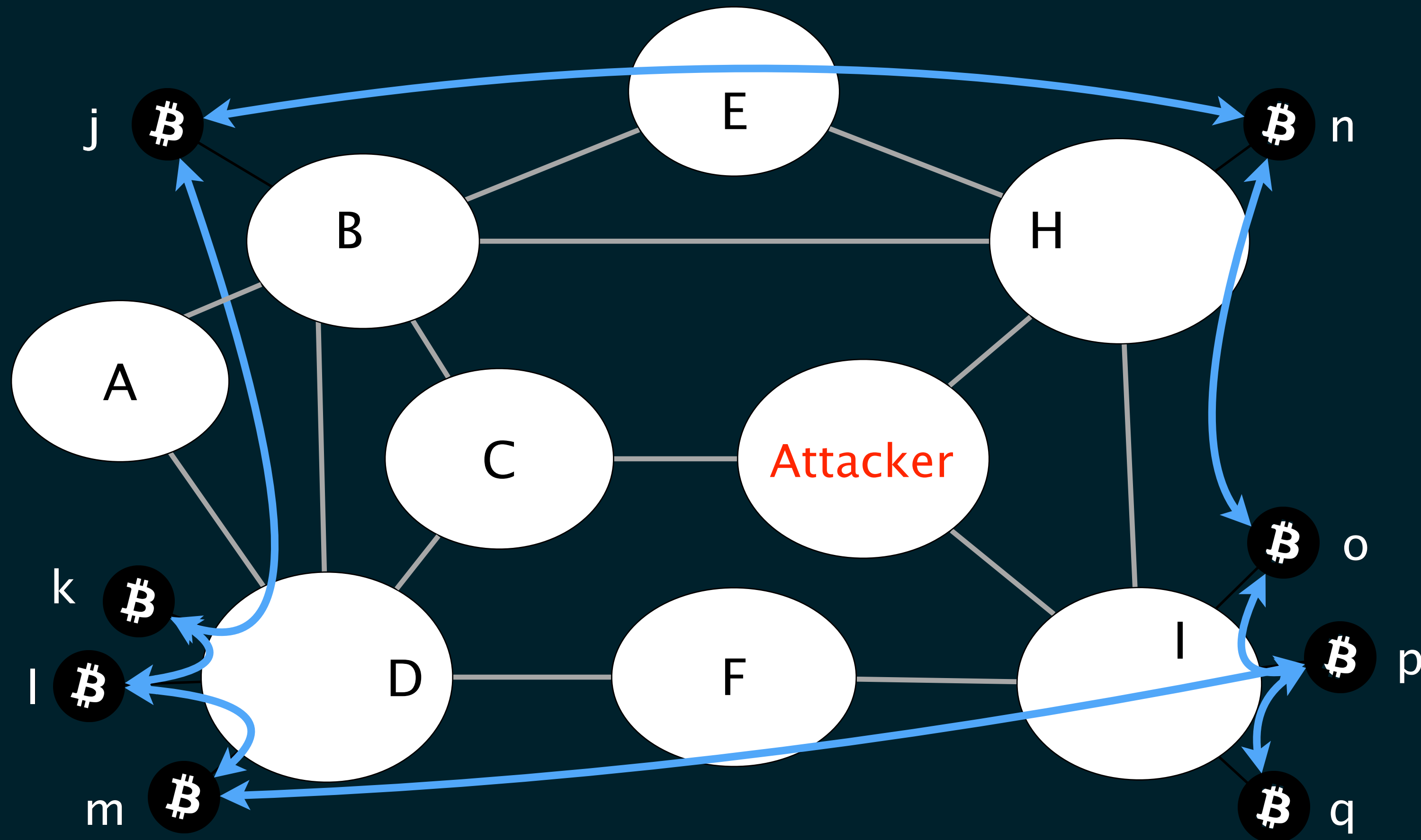




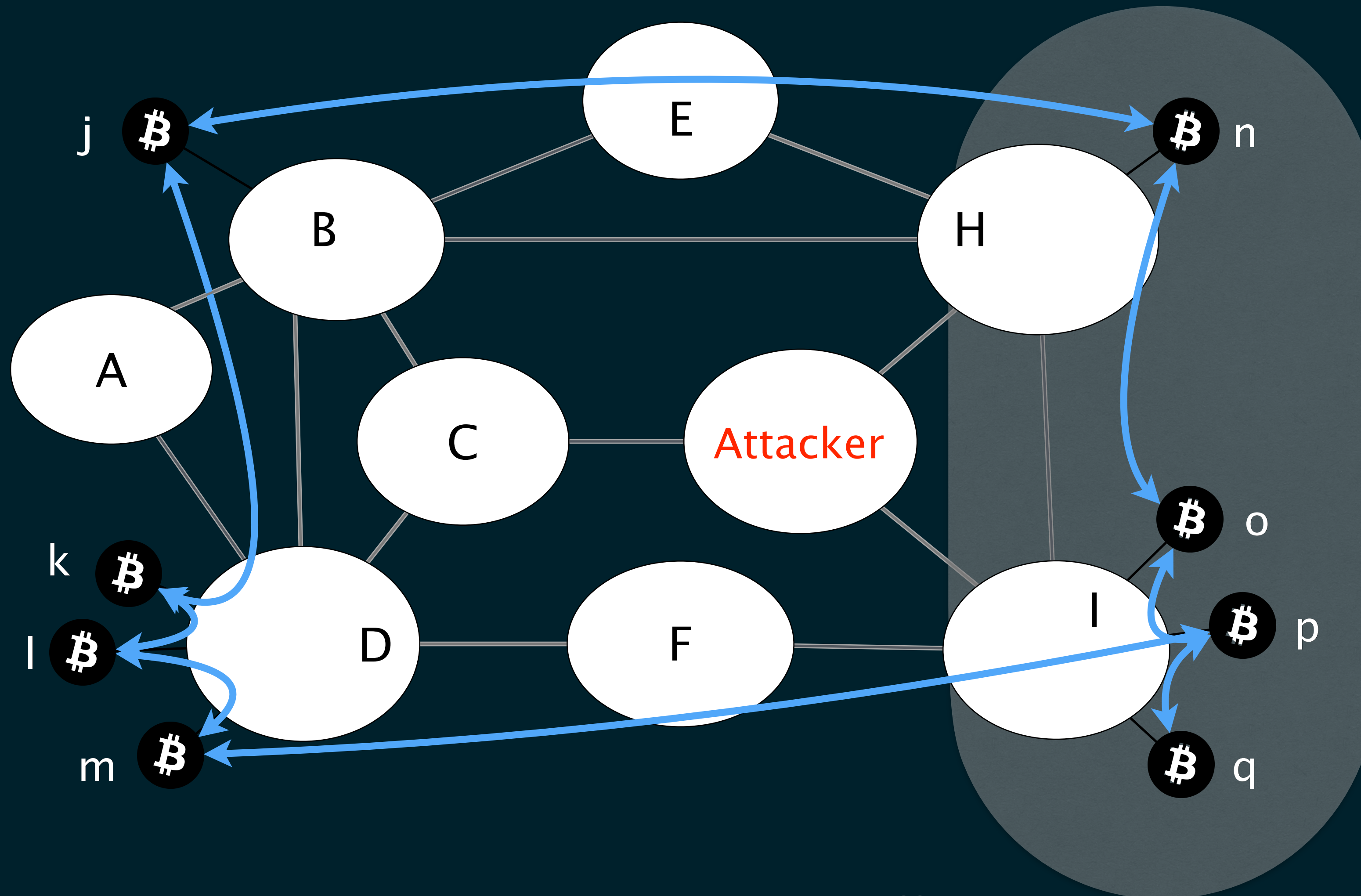
Bitcoin clients exchange Blocks  
which contain the most recent transactions



A malicious or compromised AS aims at isolating the grey zone

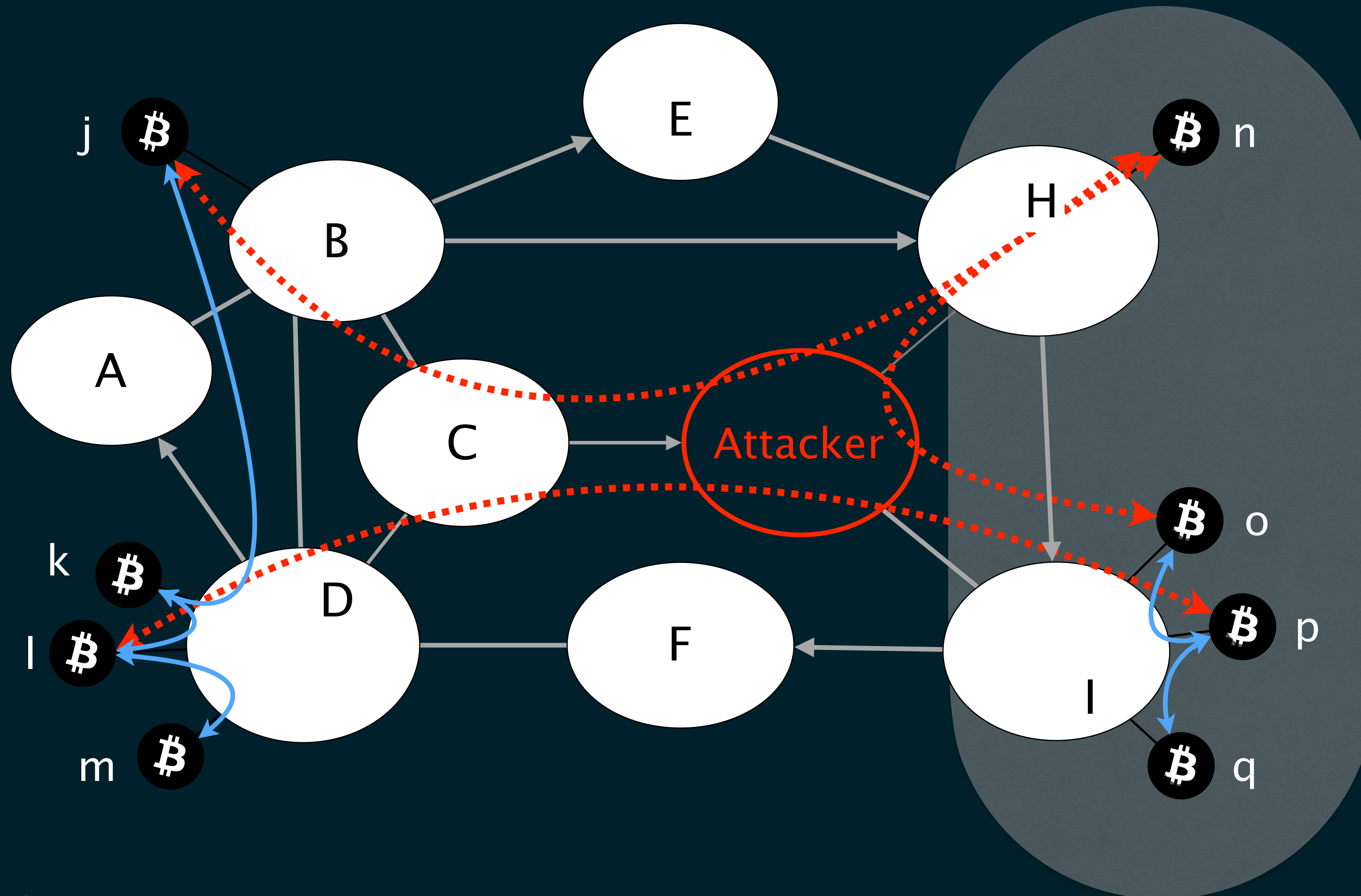


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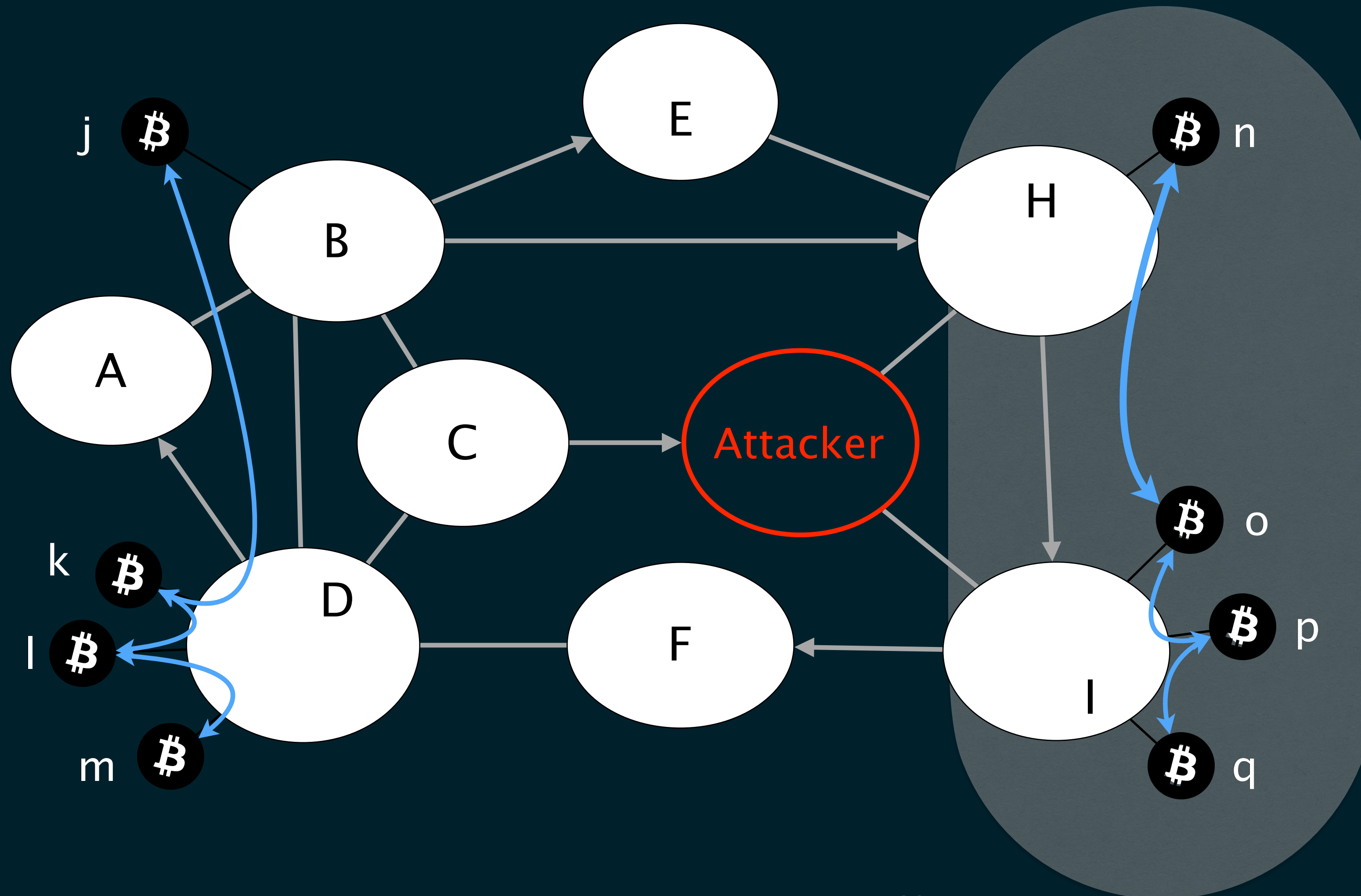




## Attacker attracts connections with BGP hijacking

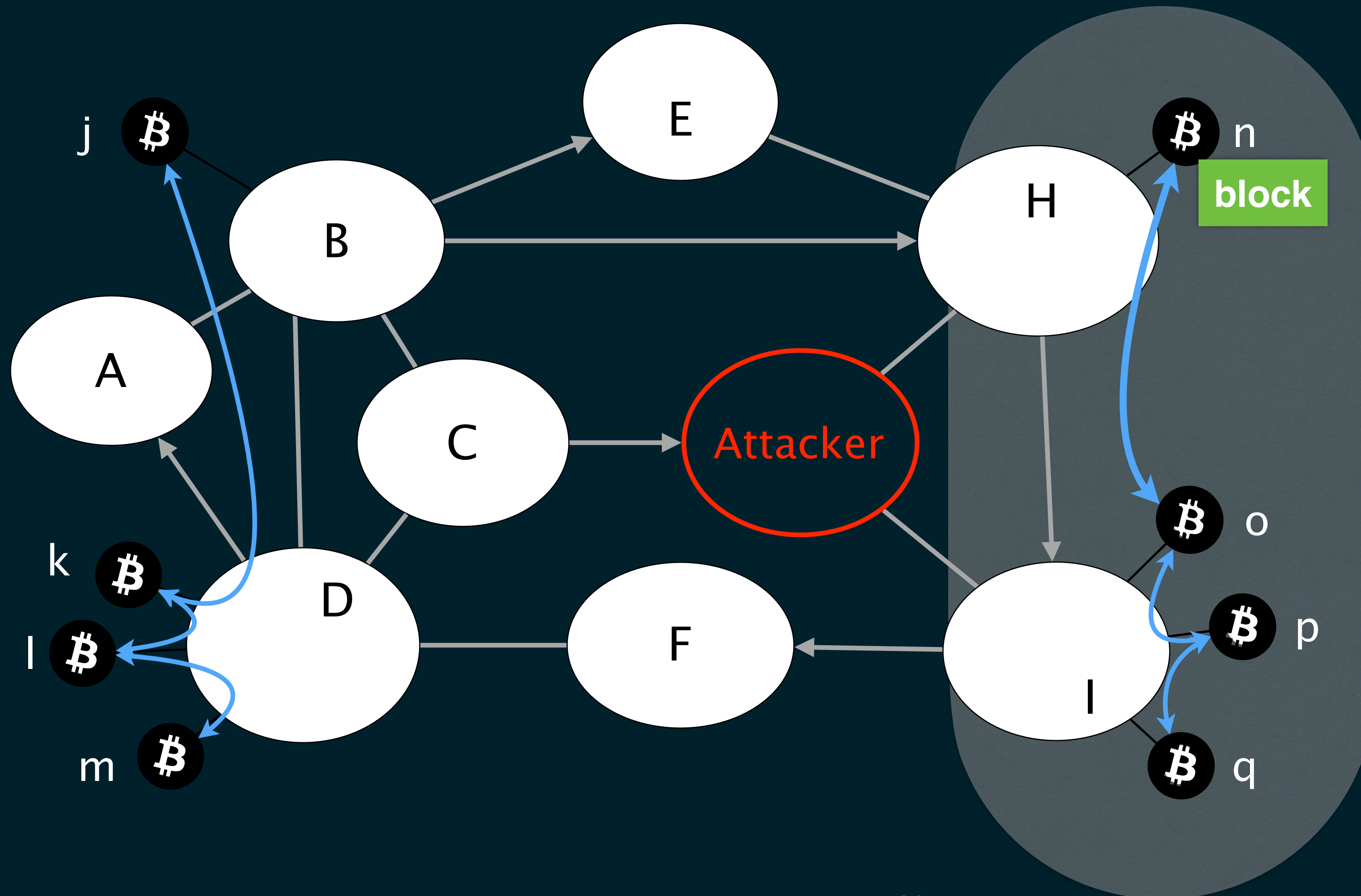


Attacker drops connections crossing the partition



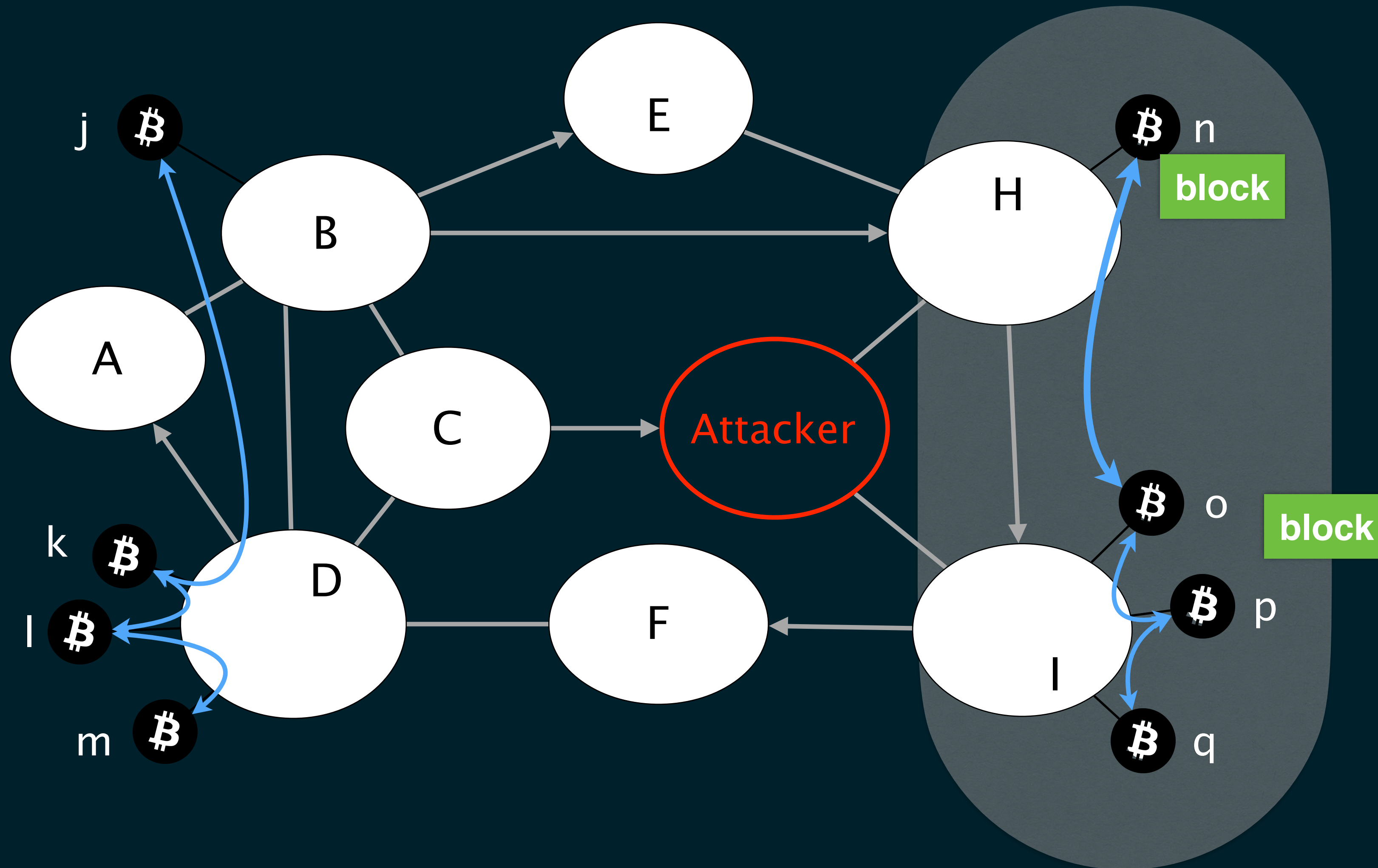


A new block in the grey zone cannot be propagated further

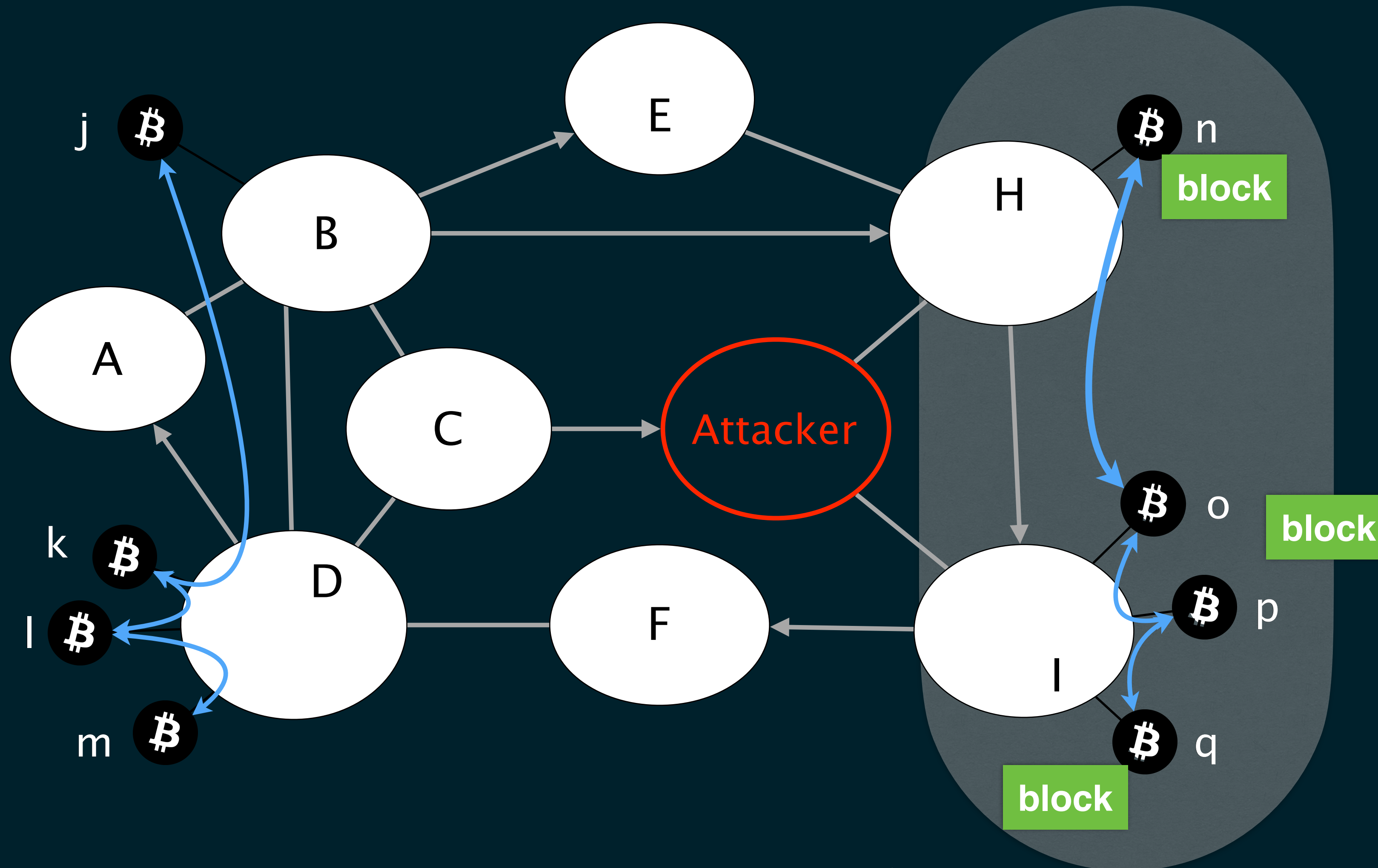




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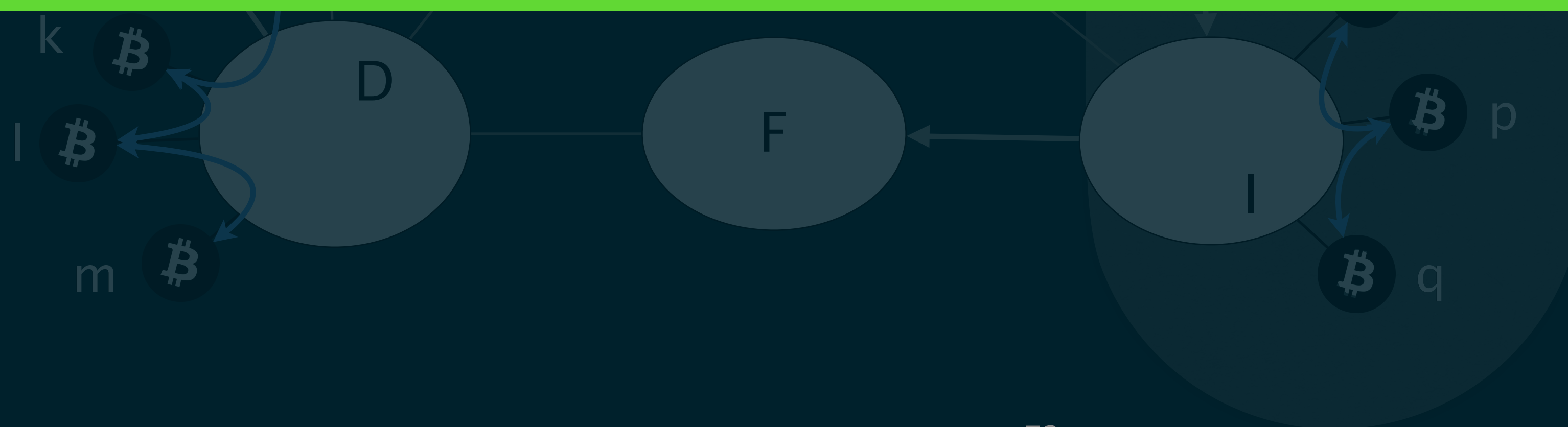


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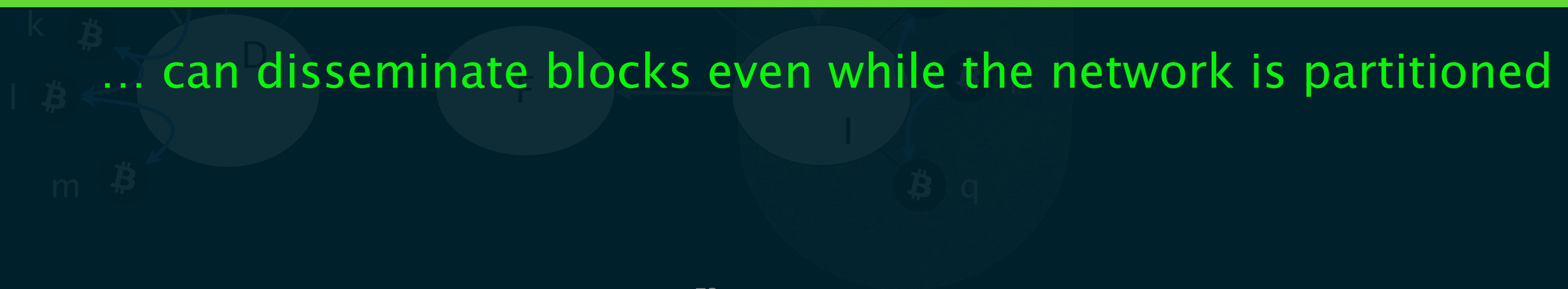
We can build an overlay of nodes strategically placed in the Internet  
s.t. they cannot be partitioned with BGP hijacks



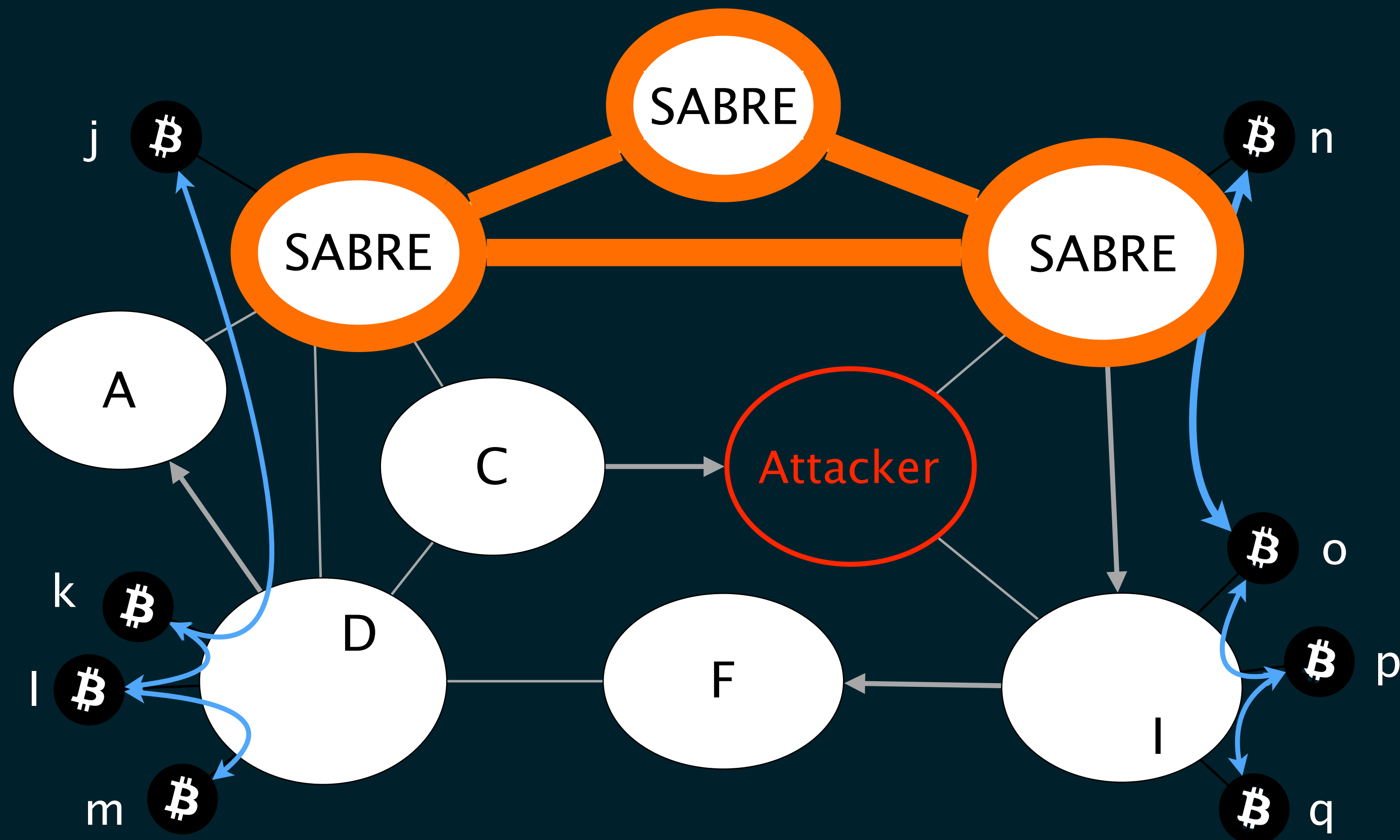


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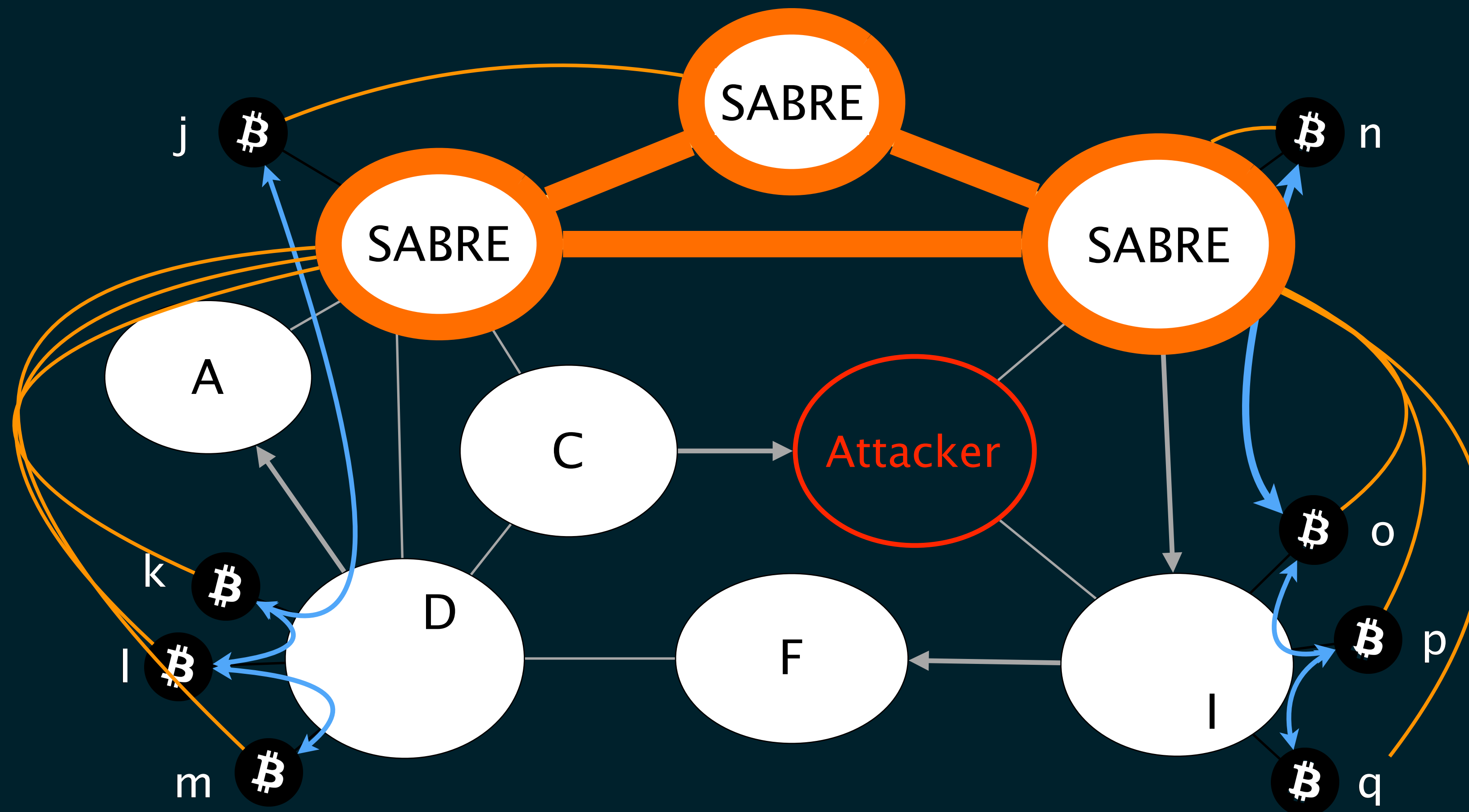
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A strongly-connected overlay can disseminate blocks even while the network is partitioned

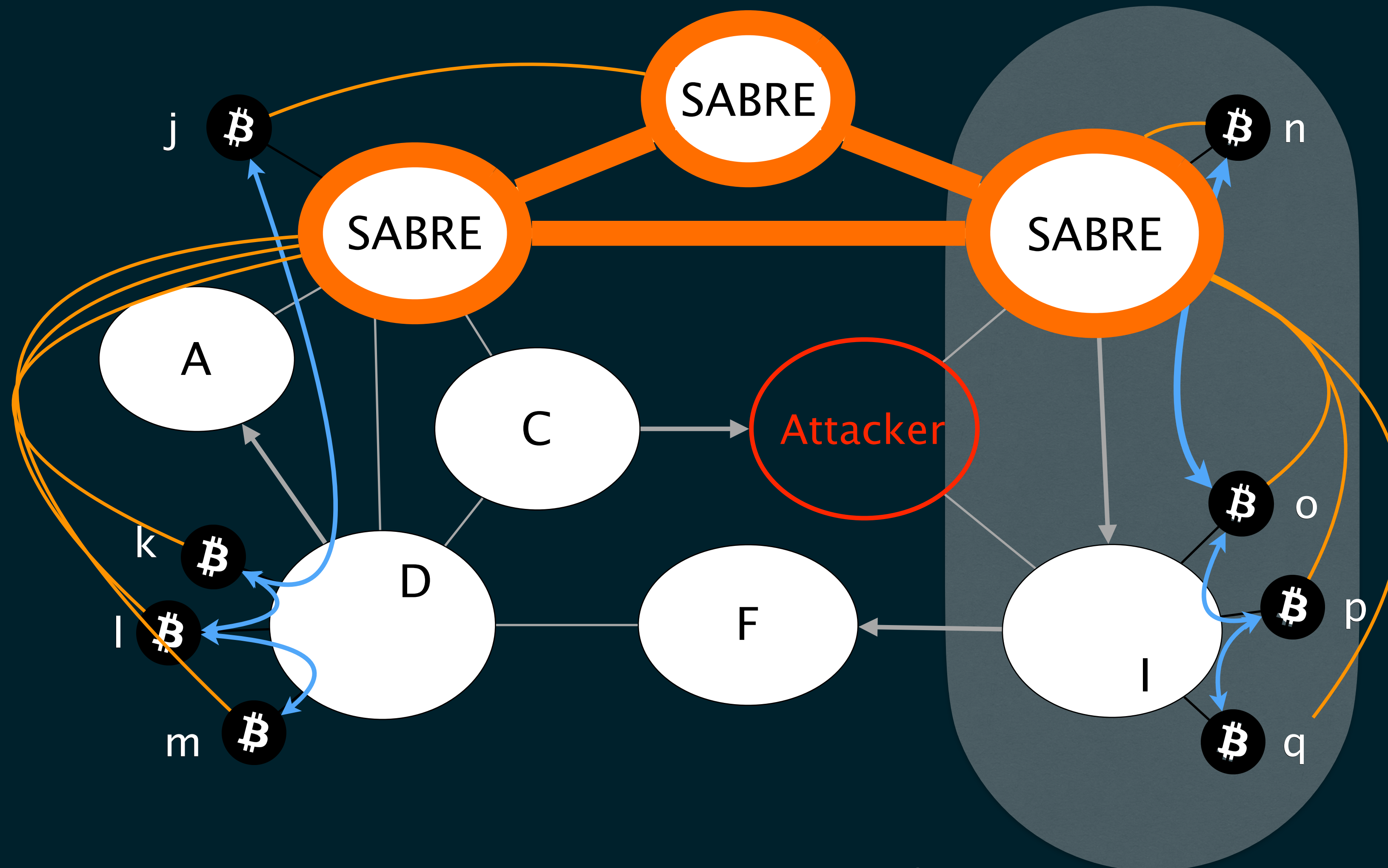


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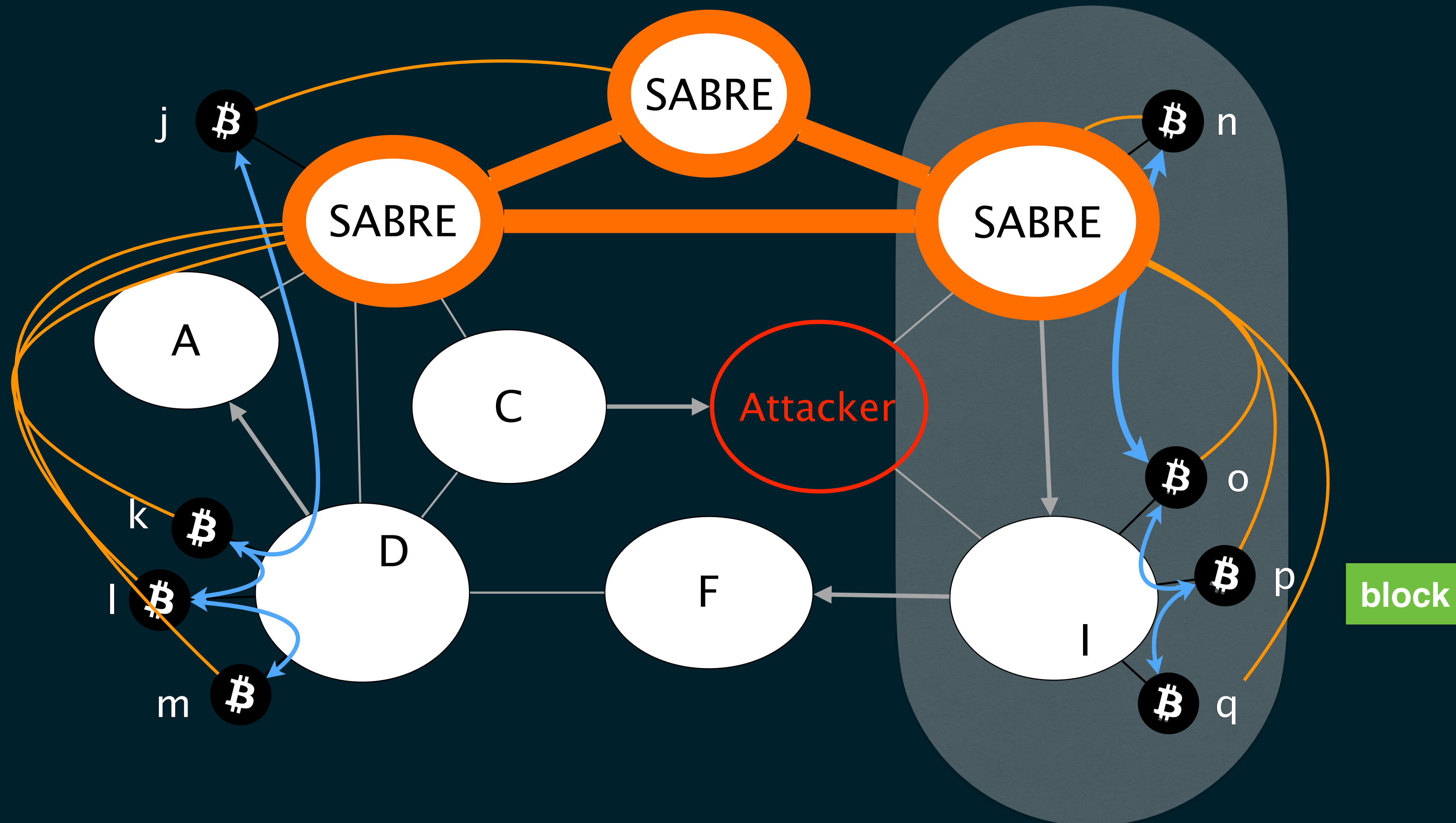




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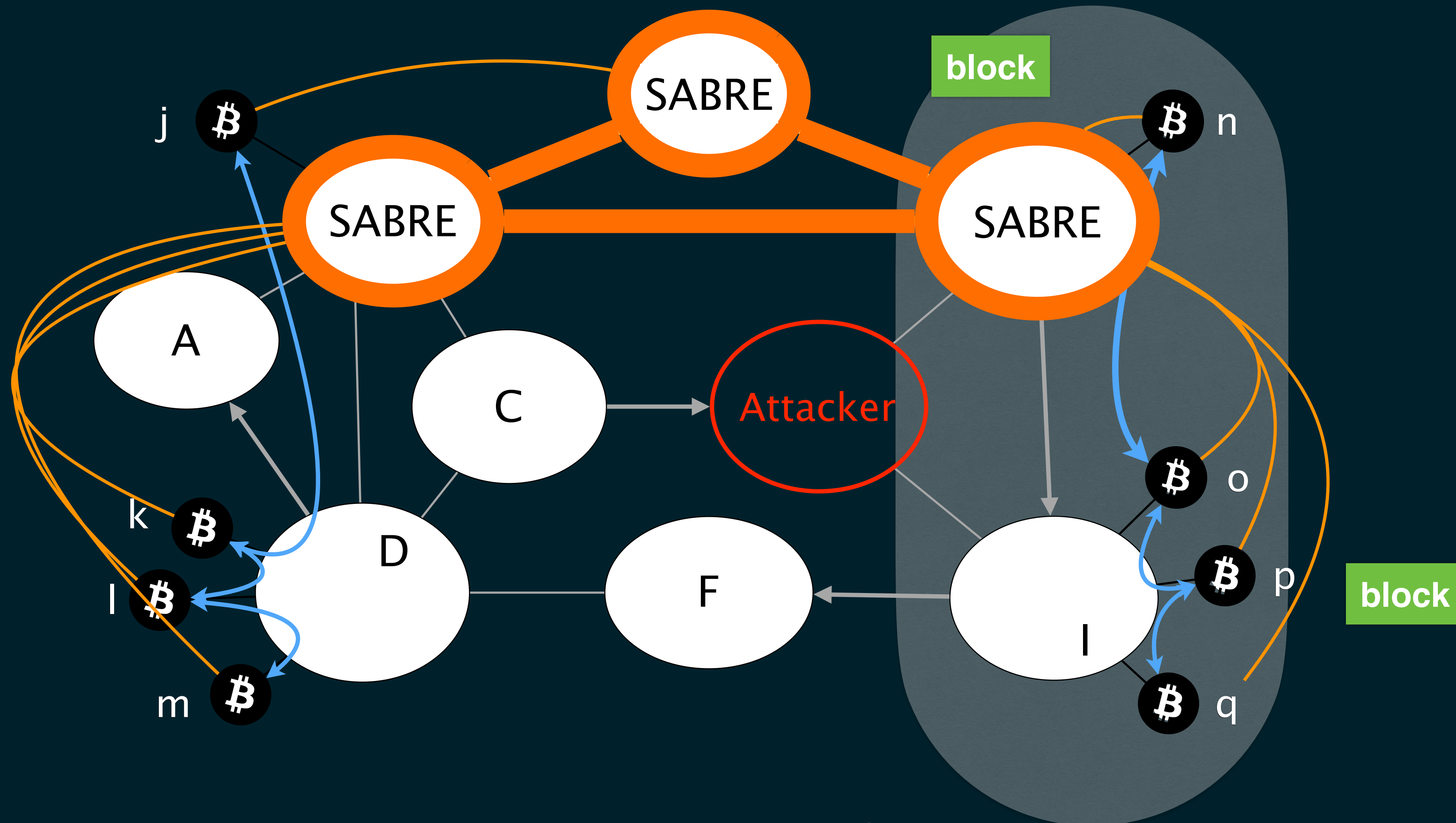


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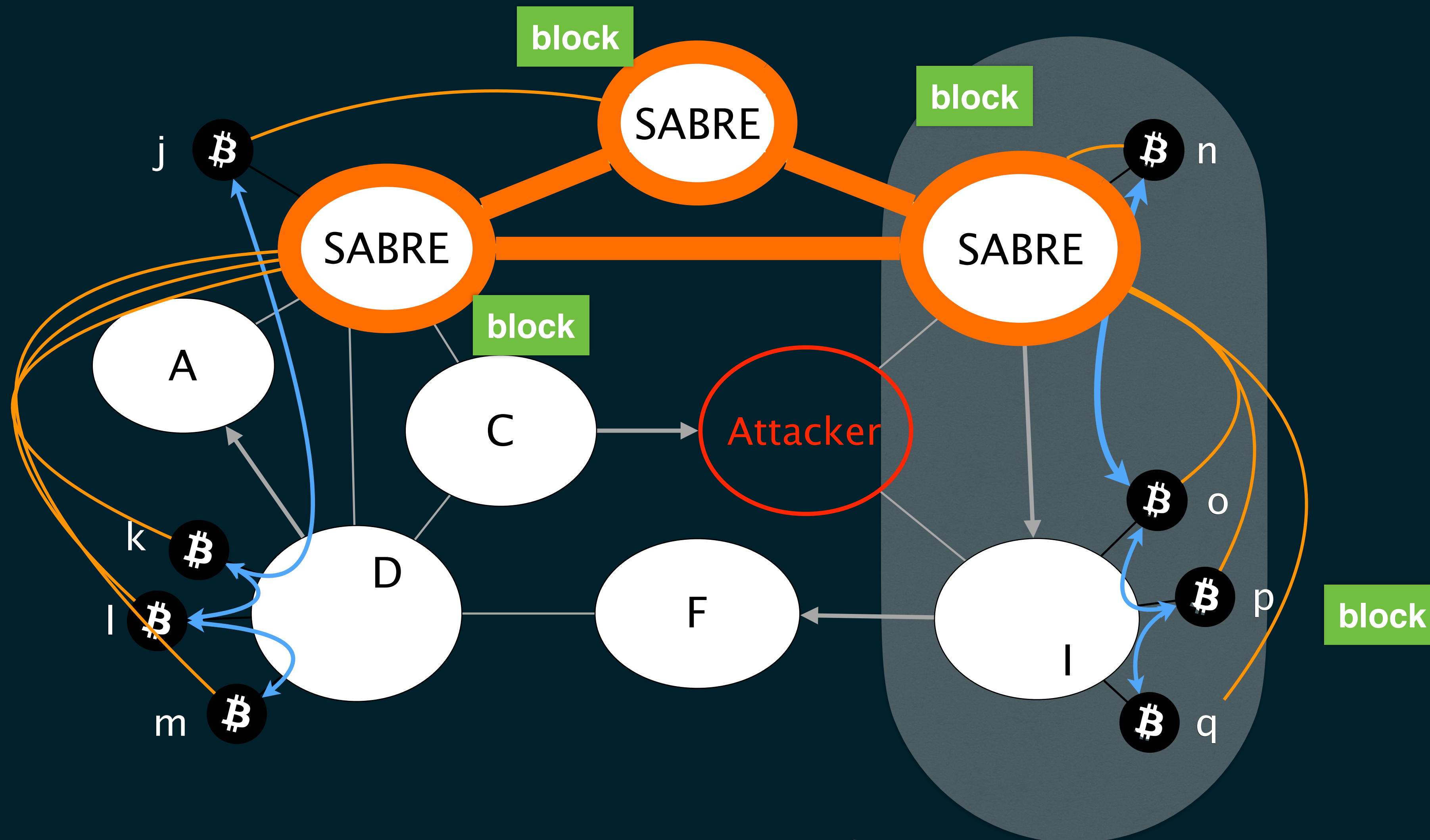


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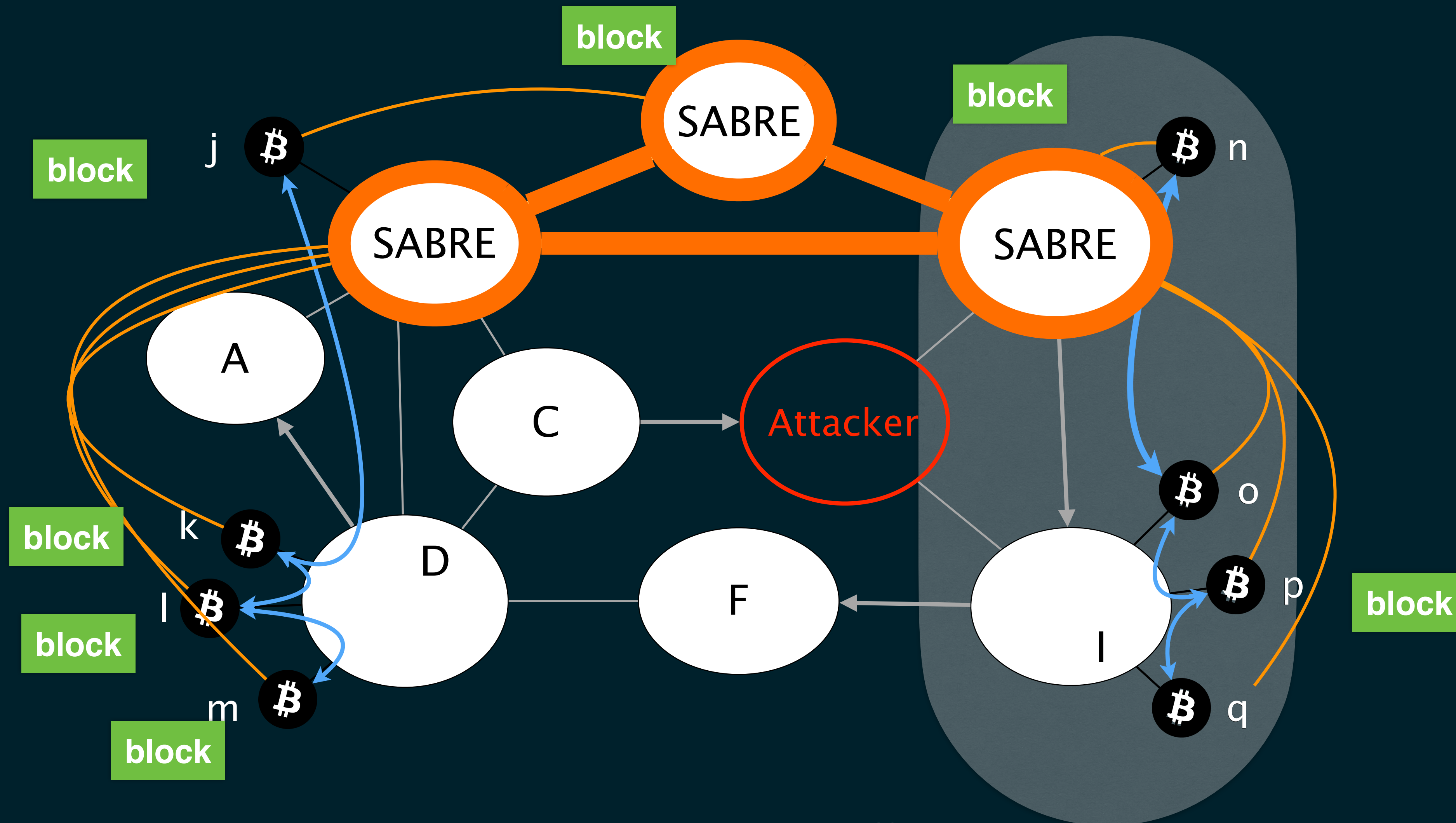




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How should the SABRE nodes be implemented?

Public SABRE nodes need to scale

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## SABRE nodes need to...

- maintain thousands of (malicious) connections
- distinguish spoofing and malicious requests
- receive, verify and relay blocks fast



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- distinguish spoofing and malicious requests
- receive, verify and relay blocks fast

*Simple software implementation would not suffice!*

SABRE can leverage programmable network devices

SABRE DP

# SABRE DP allows relay nodes to deal with high malicious or benign load

software

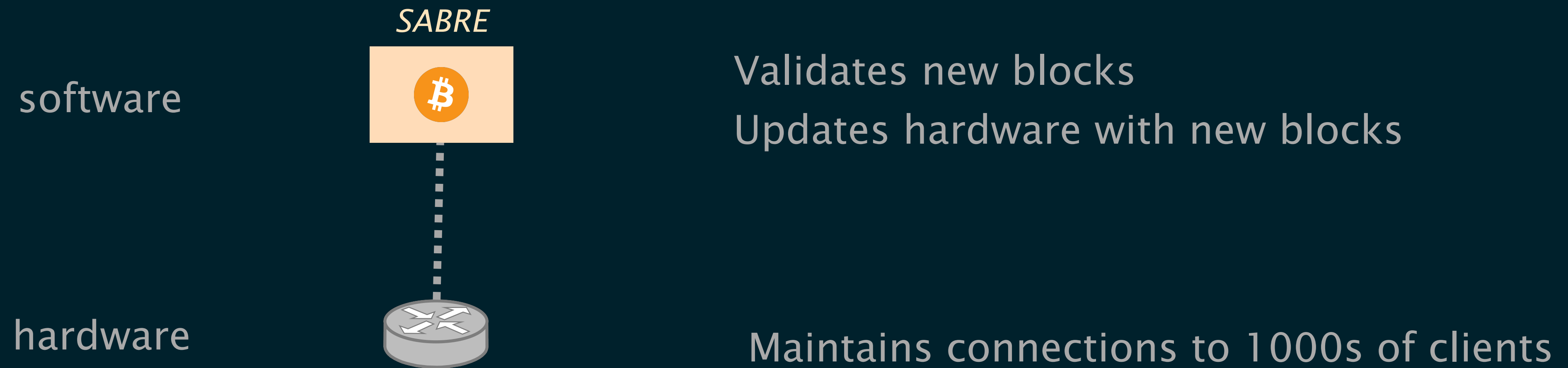
*SABRE*



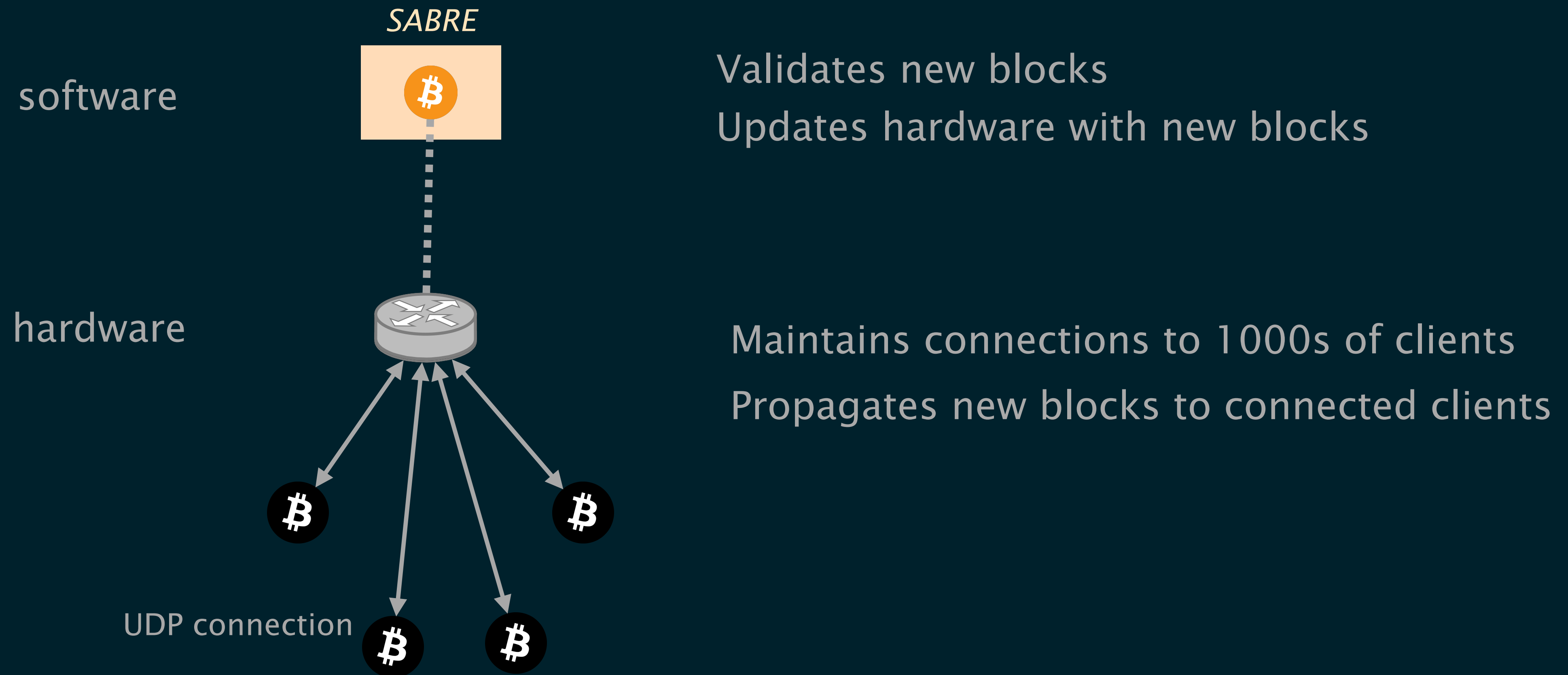
Validates new blocks  
Updates hardware with new blocks



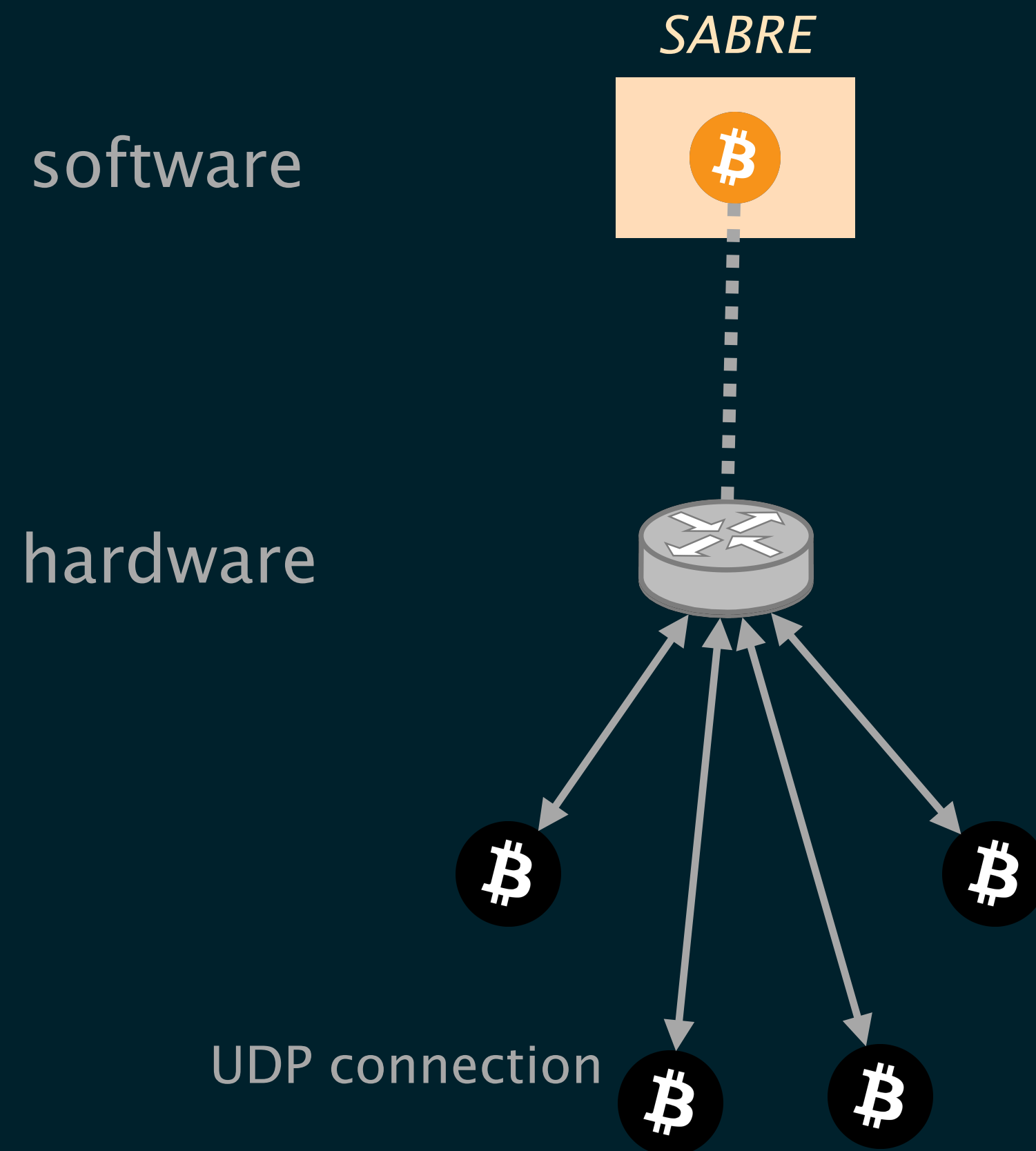
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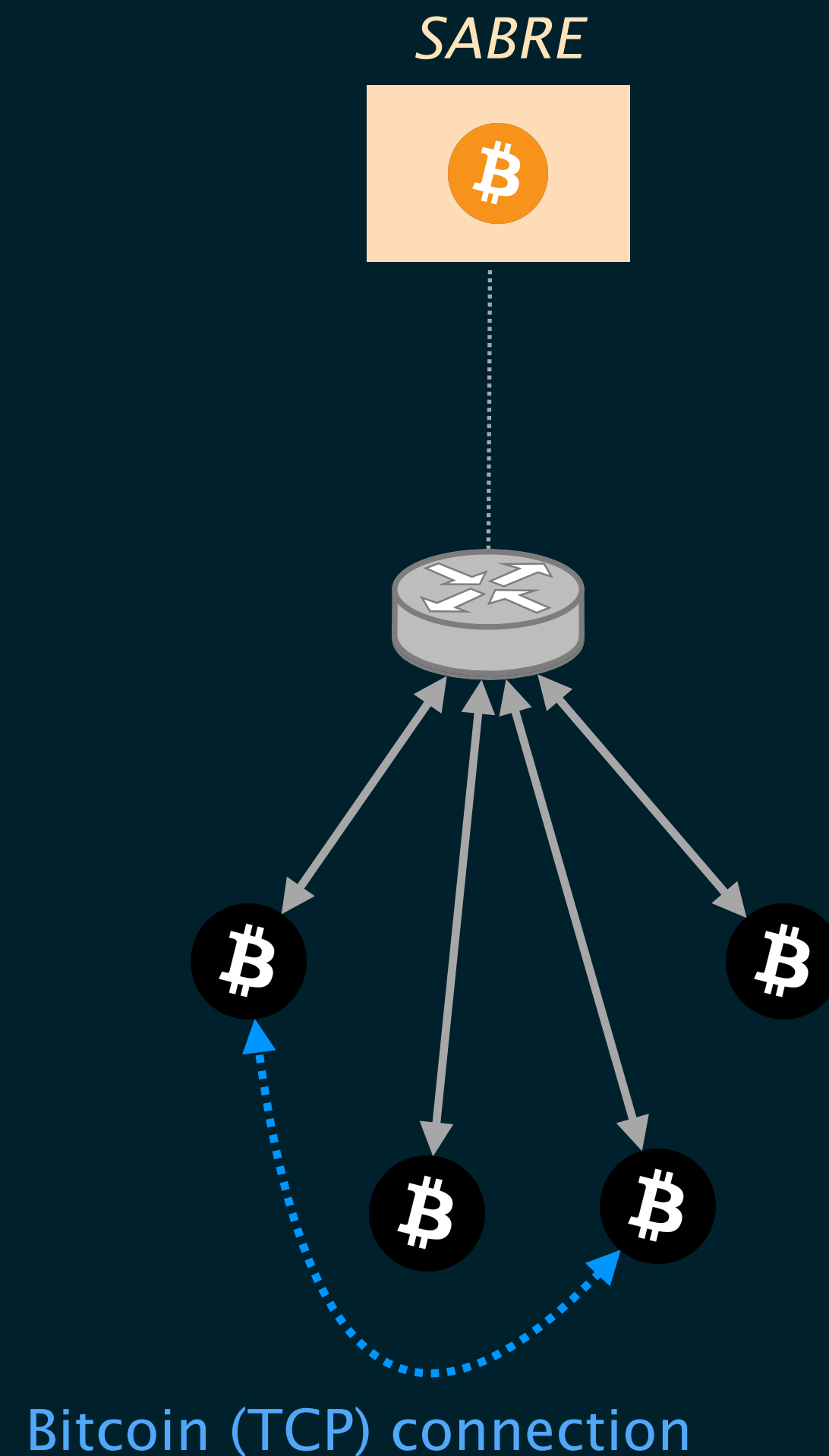


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Maintains connections to 1000s of clients  
Propagates new blocks to connected clients  
Protects the software from malicious clients



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